

Layered Coding Rule Definition and Enforcing Using LLVM

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Motivation: C++ "strange" behaviour

```
class A {  
public:  
    A();  
    virtual void func();  
};
```

```
class B : public A {  
public:  
    B() : A() {}  
    virtual void func();  
};
```

```
A::A() {  
    func();  
}
```

```
B *d = new B();
```

```
// A::func or B::func?
```

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Coding rule HICPP 3.3.13

“Do not invoke virtual methods of the declared class in a constructor or destructor.”

Definition (Coding rules)

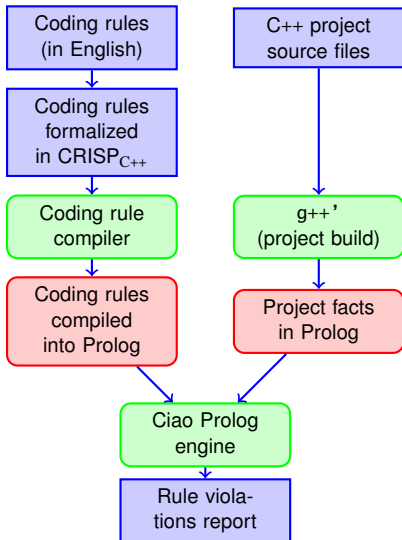
Coding Rules constrain admissible constructs of a language to help produce better code.

- Aim at improving
 - Reliability
 - Security
 - Maintainability
 - Portability
 - ...
- Coding standards
 - MISRA-C, MISRA-C++
 - Parasoft's High-Integrity C++ (HICPP)
 - CERT's (CMU) Secure Coding Standards for C, C++, Java
- Need of automation and extensibility
- Problem: big gap between semi-formal definition of rules and actual implementation

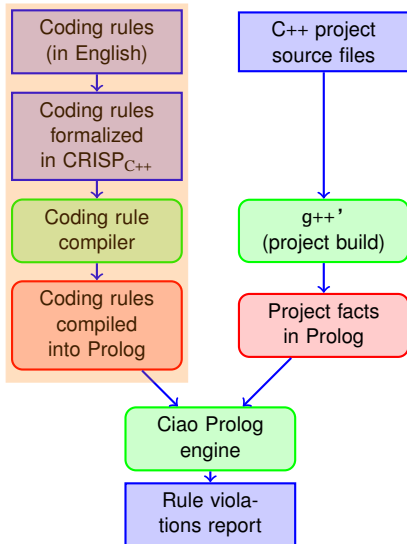
Diversity of Coding Rules

- "Do not use the 'inline' keyword for member functions."
- "Do not call the malloc() function."
- "Expressions that are effectively Boolean should not be used as operands to operators other than (&&, || and !)."
- "If a virtual function in a base class is not overridden in any derived class, then make it non virtual."
- "All automatic variables shall have been assigned a value before being used."
- "Behaviour should be implemented by only one member function in a class."

Previous prototype based on GCC

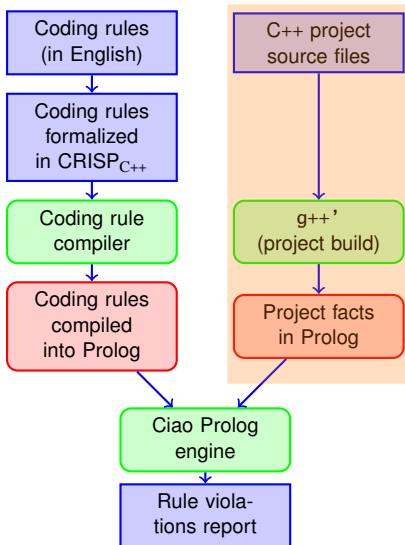


Previous prototype based on GCC



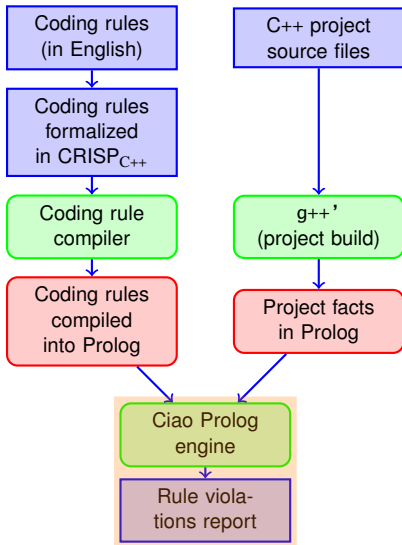
- 1 Coding rule(s) written **once** in a logic-based formalism called CRISP
- 2 Extract program information using GCC, and store it
- 3 Search (using a Prolog engine) for a counterexample

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Example of rule formalisation in CRISP

Rule HICPP 3.3.13

“Do not invoke virtual methods of the declared class in a constructor or destructor.”

Example of rule formalisation in CRISP

Rule HICPP 3.3.13

“Do not invoke virtual methods of the declared class in a constructor or destructor.”

```
rule          HICPP 3.3.13
violated by  Caller : MemberFunction; Callee : VirtualFunction
when        exists R : Record such that (
            R hasMember Caller
            and R hasMember Callee
            and (
                Caller is Constructor
                or Caller is Destructor
            )
            and Caller calls+ Callee
        )
.
```

- Sorts

*Variable, DataMember, LocalVariable
 Function, MemberFunction, Constructor
 Type, PointerType, Record
 Scope, Namespace, Record, CompoundStatement
 RecordMember*

- Predicates

Function	<i>calls</i>	Function
Record	<i>hasImmediateBase</i>	Record
Variable	<i>hasType</i>	NonFunctionType
Function	<i>hasType</i>	FunctionType
Function	<i>isDefinedIn</i>	Scope
Record	<i>hasMember</i>	RecordMember
PointerType	<i>hasPointedType</i>	Type
FunctionType	<i>hasReturnType</i>	Type

Rules that need some kind of static analysis

Rule HICPP 3.4.2

“Do not return non-const handles to class data from *const* member functions.”

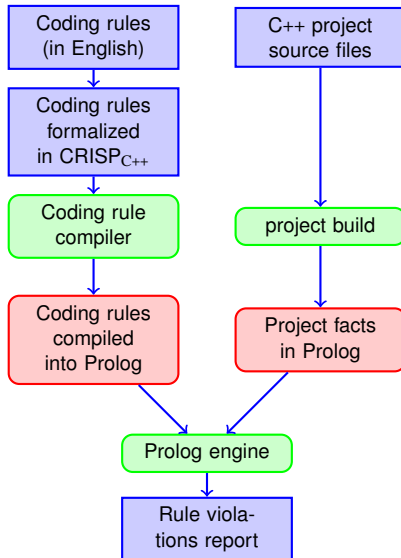
```
class A {
public:
    int* foo() const {
        return m_pa;    // permits subsequent mod. of private data
    }

private:
    int* m_pa;
};

void bar() {
    const A a;
    int* pa = a.foo();
    *pa = 10;          // modifies private data in a!
};
```

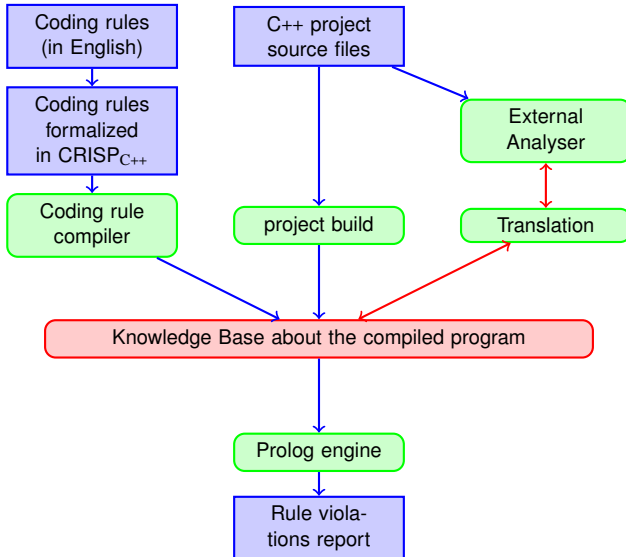
Layered rule definition

Integration of information from external analysers



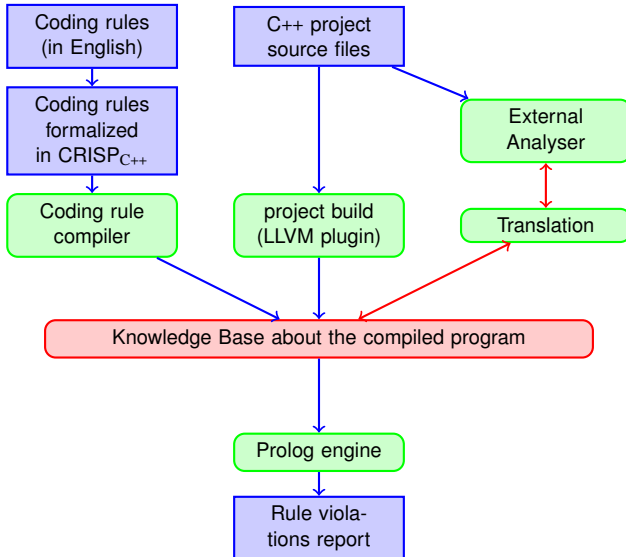
Layered rule definition

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Integration of information from external analysers



LLVM: Low-Level Virtual Machine

- "Modular and reusable compiler and toolchain technologies" (<http://llvm.org>)
- Useful components and features
 - clang C/C++ compiler front-end
 - Well specified code representation (LLVM IR)
 - Target-independent optimiser/analyser (`-analyze`)
 - Alias analysis
 - Link-time optimisation (and analysis)
- LLVM IR is front-end independent
 - Advantage: analysis info can be reused across languages
 - Disadvantage: need to keep track of the correspondence between source-level constructions and IR elements

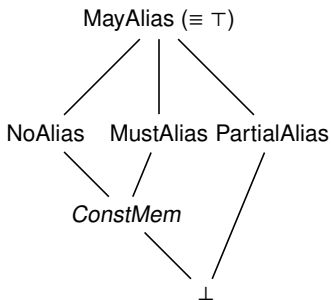
- API

```
enum AliasResult { NoAlias      = 0 , MayAlias   = 1  
                  , PartialAlias = 2 , MustAlias  = 3 }
```

```
virtual AliasResult  
alias (const Location &LocA, const Location &LocB);
```

- A **Location** represents a memory location in the LLVM IR (assembler) code: Pointer + Size
- Diversity of implementations
 - -basicaa
 - -steens-aa: Steensgaard's algorithm
 - -ds-aa: Data Structure Analysis algorithm
 - -scev-aa: Scalar Evolution queries
 - -tbaa: Type-based

A domain for alias analysis results



- Complementary LLVM API

virtual bool pointsToConstantMemory (const Location &Loc);

- CRISP expressions for alias queries, examples:

L1 *alias* L2 \sqsubseteq NoAlias

L1 *alias* L2 = PartialAlias

- Idea: use the above expressions as constraints in CLP

LLVM IR for counterexample of Rule HICPP 3.4.2

```
class A {
public:
    int* foo(bool b) const {
        int* local = 0;
        if (b) local = m_pa;
        return local;
    }
private:
    int* m_pa;
};

define linkonce_odr i32* @_ZNK1A3fooEb(%class.A* nocapture %this, i1 zeroext %b)
    nounwind readonly align 2 {
entry:
    tail call void @llvm.dbg.value(metadata !{%class.A* %this}, i64 0, metadata !28)
    tail call void @llvm.dbg.value(metadata !42, i64 0, metadata !30)
    br i1 %b, label %if.then, label %if.end

if.then:                                     ; preds = %entry
    %m_pa = getelementptr inbounds %class.A* %this, i64 0, i32 0
    %tmp2 = load i32** %m_pa, align 8
    tail call void @llvm.dbg.value(metadata !{i32* %tmp2}, i64 0, metadata !30)
    br label %if.end

if.end:                                       ; preds = %if.then, %entry
    %local.0 = phi i32* [ %tmp2, %if.then ], [ null, %entry ]
    ret i32* %local.0
}
```

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class A {
public:
  int* foo(bool b) const {
    int* local = 0;
    if (b) local = m_pa;
    return local;
  }
private:
  int* m_pa;
};
```

- New CRISP predicates

MemberFunction	<i>hasThisLocation</i>	Location
Function	<i>returnsLocation</i>	Location
Location	<i>isGottenAsOffsetFrom</i>	Location

```
define linkonce_odr i32* @_ZNK1A3fooEb(%class.A* nocapture %this, i1 zeroext %b)
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entry:
  tail call void @llvm.dbg.value(metadata !{%class.A* %this}, i64 0, metadata !28)
  tail call void @llvm.dbg.value(metadata !42, i64 0, metadata !30)
  br i1 %b, label %if.then, label %if.end

if.then:
  ; preds = %entry
  %m_pa = getelementptr inbounds %class.A* %this, i64 0, i32 0
  %tmp2 = load i32** %m_pa, align 8
  tail call void @llvm.dbg.value(metadata !{i32* %tmp2}, i64 0, metadata !30)
  br label %if.end

if.end:
  ; preds = %if.then, %entry
  %local.0 = phi i32* [ %tmp2, %if.then ], [ null, %entry ]
  ret i32* %local.0
}
```

CRISP definition of Rule HICPP 3.4.2

Rule HICPP 3.4.2

“Do not return non-const handles to class data from const member functions.”

CRISP definition of Rule HICPP 3.4.2

Rule HICPP 3.4.2

“Do not return non-const handles to class data from const member functions.”

```
rule          HICPP 3.4.2
violated by  M : ConstMemberFunction
when        exists Data, This, Ret : Location;
            FuncT, RetT, PointT : Type
            such that (
                M hasType T
                and T hasReturnType RetT
                and RetT is PointerType
                and RetT hasPointedType PointT
                and not PointT is ConstType
                and M hasThisLocation This
                and M returnsLocation Ret
                and Data isGottenAsOffsetFrom This
                and Ret alias Data  $\not\sqsubseteq$  NoAlias
            )
```

- Define more rules that need Alias Analysis and find the right abstractions of low-level concepts.
- Find other analyses useful for formalising coding rules and define suitable domains for them.
- Devise a method for solving analysis domain constraints.
- Study rules in which results from different analyses are necessary, and how to combine them.