

Fuzzy Prolog

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Overview

- Basics

- Introduction
- Description
- Implementation

- Extensions

- Incompleteness
- Constructive negative Queries
- Discrete Fuzzy Sets
- Collaborative Fuzzy Agents
- Fuzzy Rules with Credibility: RFuzzy

- Work proposals

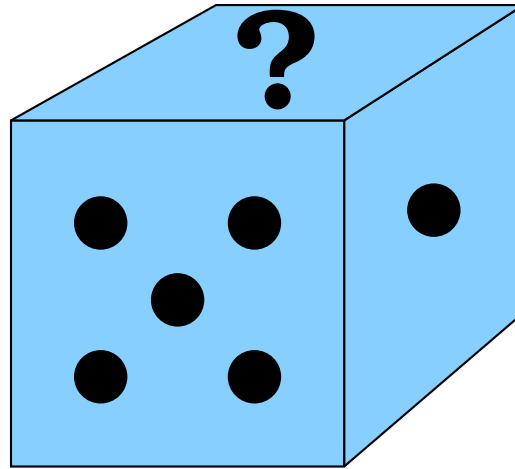
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Modeling Real World

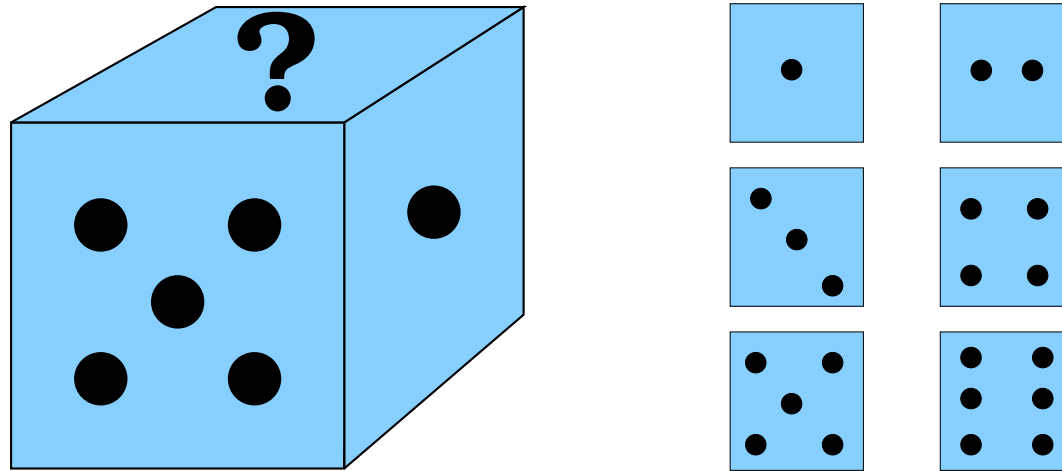
- Knowledge:
 - Uncertainty
 - Probability
 - Fuzziness
 - Incompleteness
 - Distributed knowledge
- Reasoning:
 - Logic

Uncertainty-Probability-Fuzziness



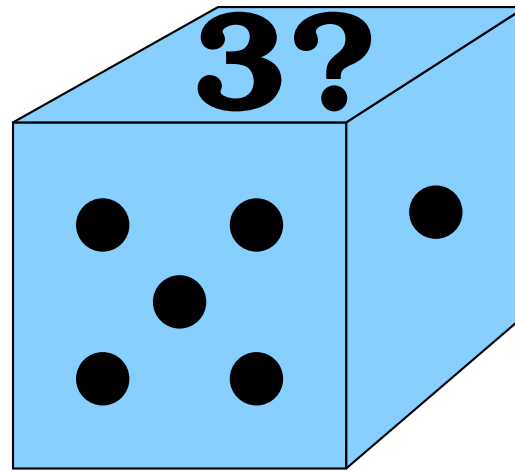
- If we throw the cube... which value will appear at the top?

Uncertainty



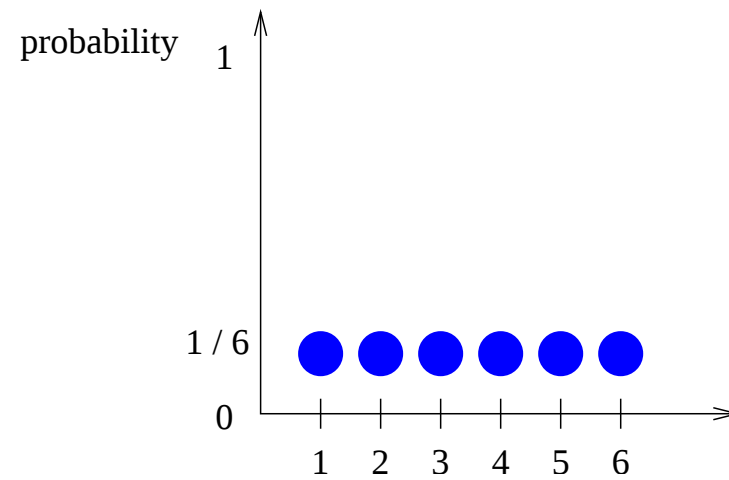
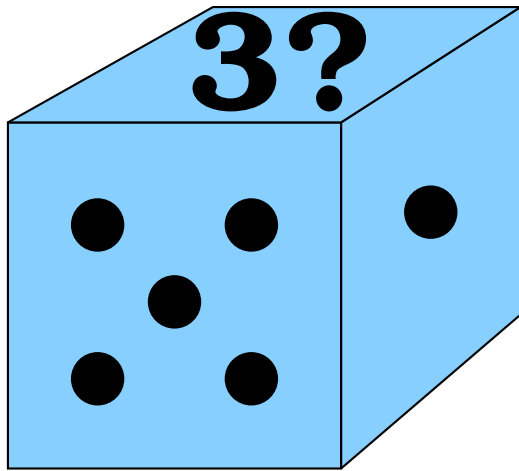
$$X = 1 \vee X = 2 \vee X = 3 \vee X = 4 \vee X = 5 \vee X = 6$$

Uncertainty-Probability-Fuzziness



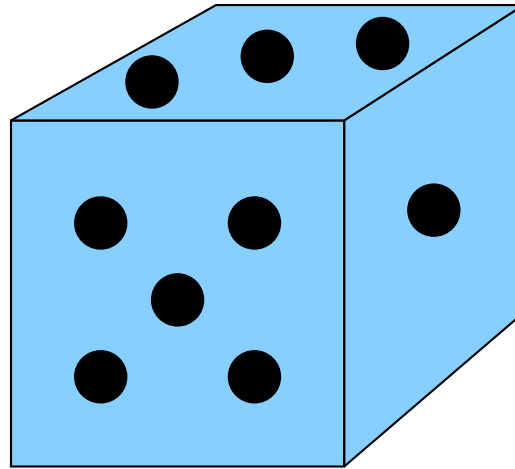
- If we throw the cube... is it probable to obtain 3 at the top?

Probability



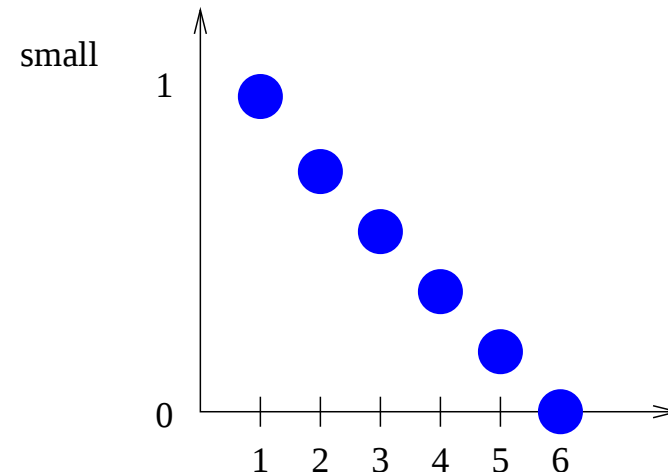
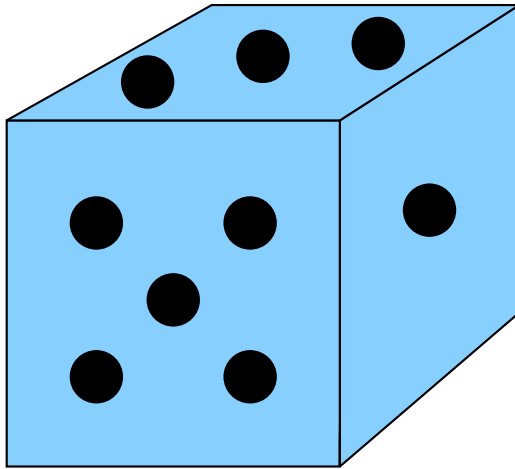
$$Pr(X = 3) = 1/6 = 0.16$$

Uncertainty-Probability-Fuzziness



- If we obtain 3 at the top... is it a small value?

Fuzziness



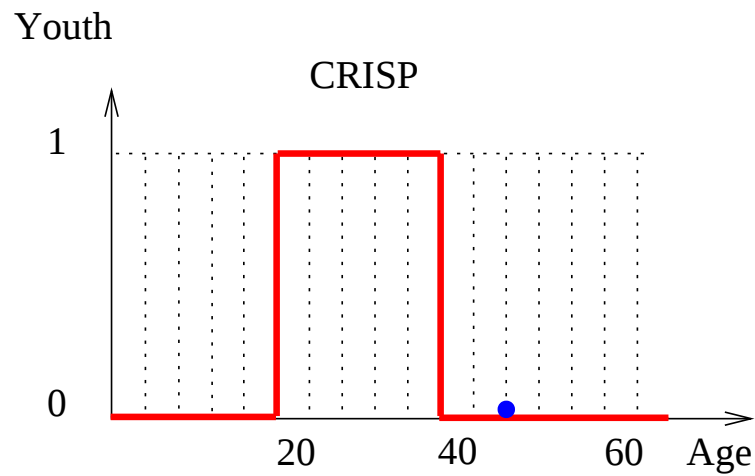
$$\textit{small}(X = 3) = 0.6$$

Value 3 is slightly small

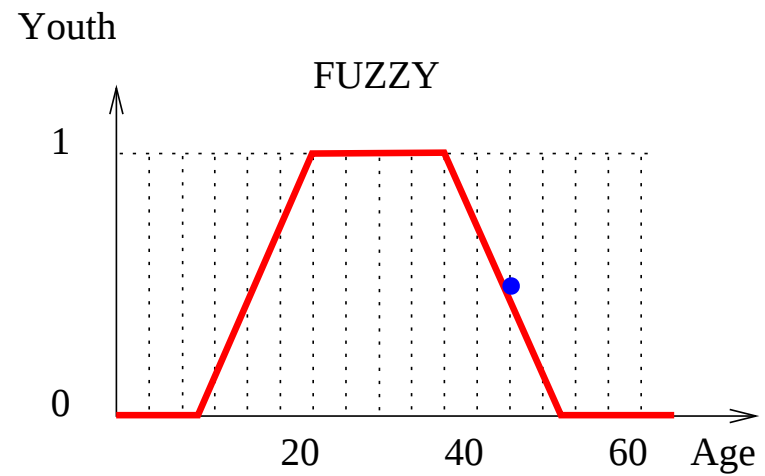
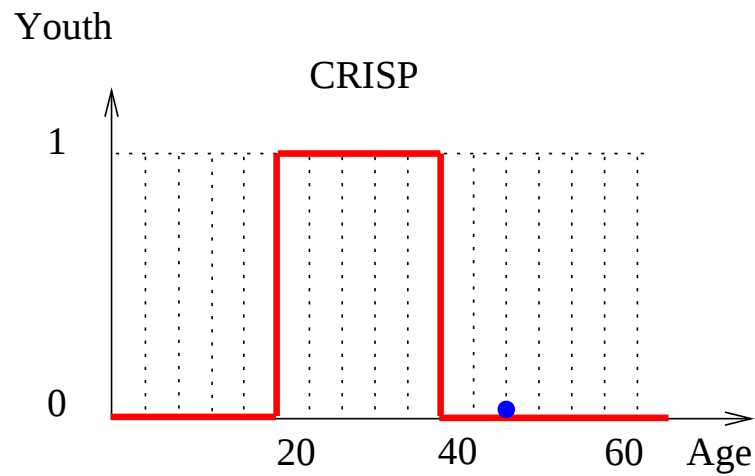
Fuzziness level: Example

Let's define the concept of youth

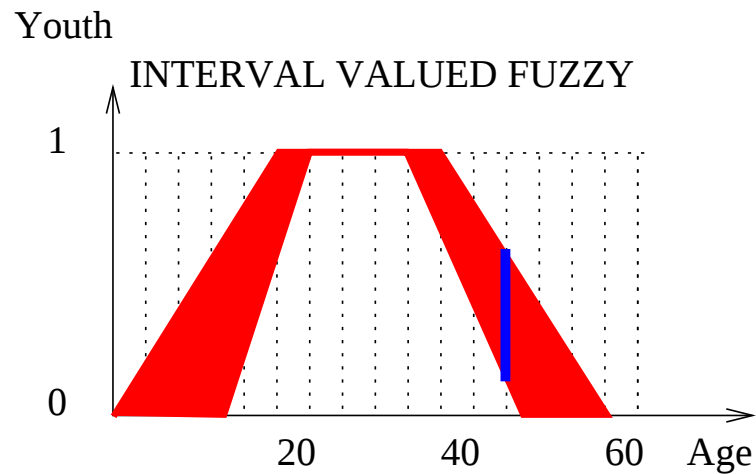
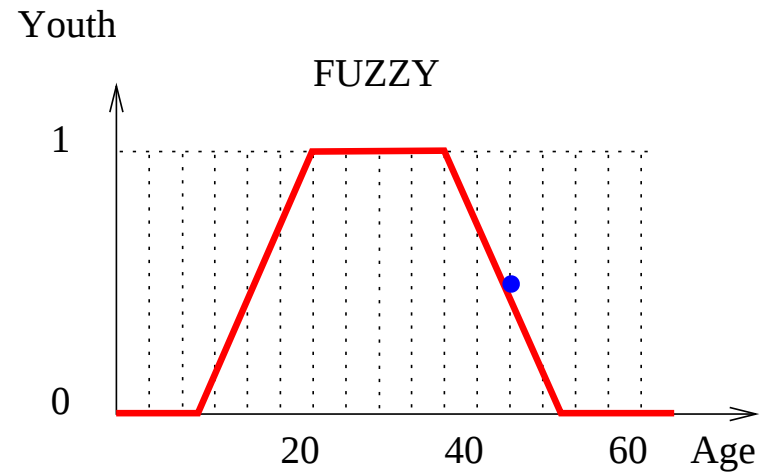
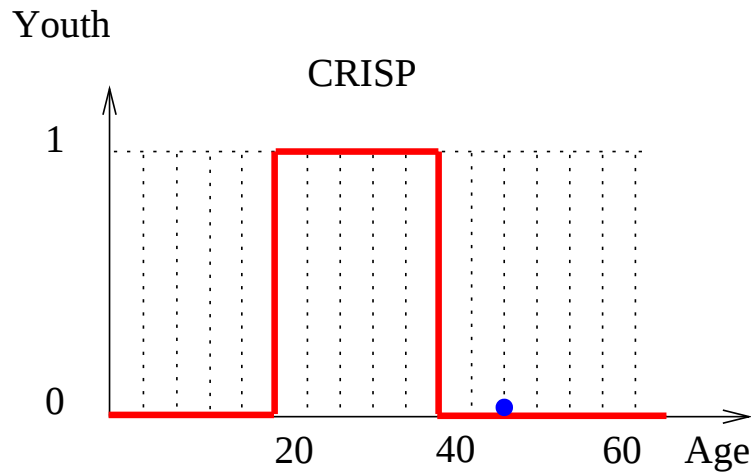
Fuzziness level



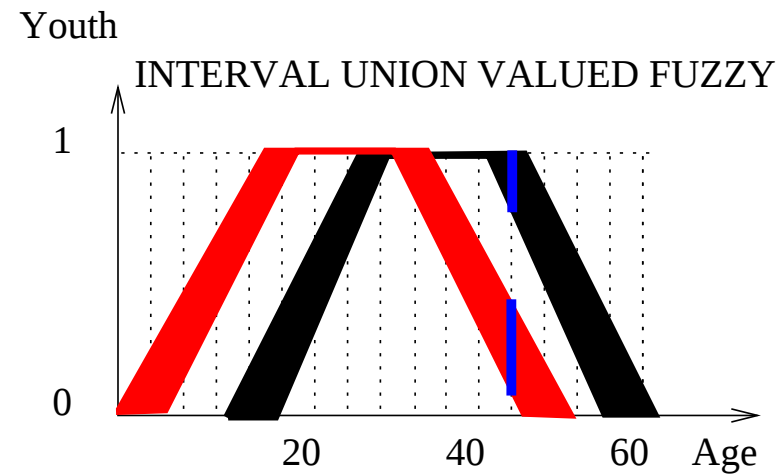
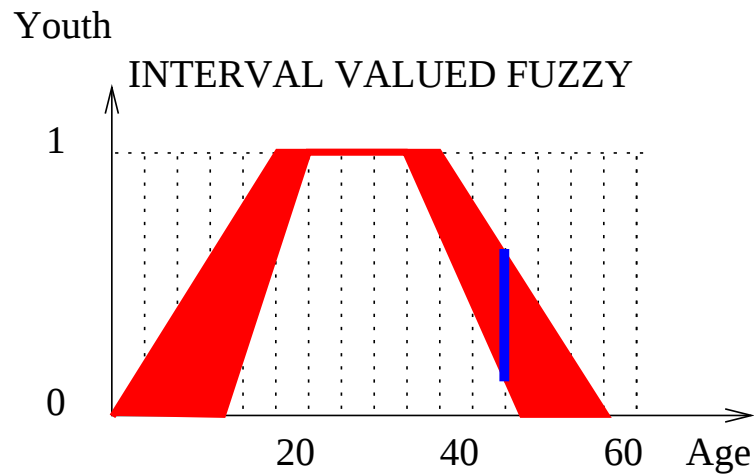
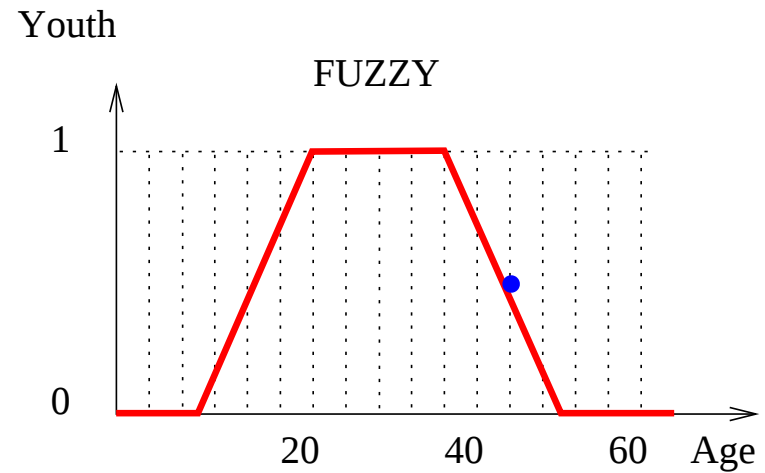
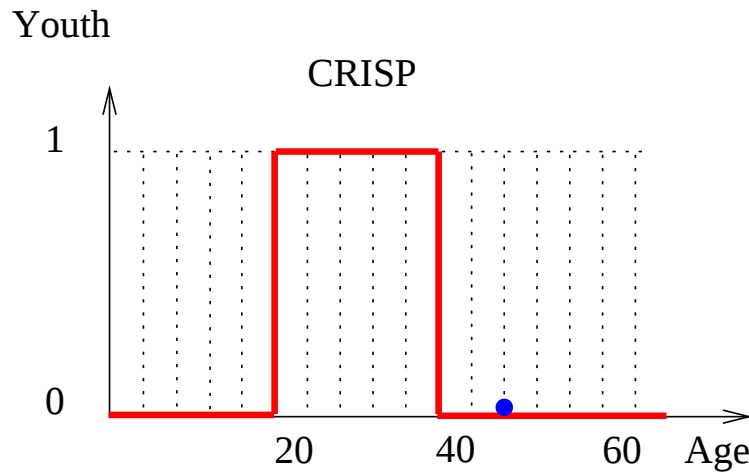
Fuzziness level



Fuzziness level



Fuzziness level



Truth value (Fuzziness level)

The value of youth of a 42 years-old man

- $V = 0$
- $V = 0.5$
- $V \in [0.2, 0.6]$
- $V \in [0.2, 0.5] \cup [0.8, 1]$

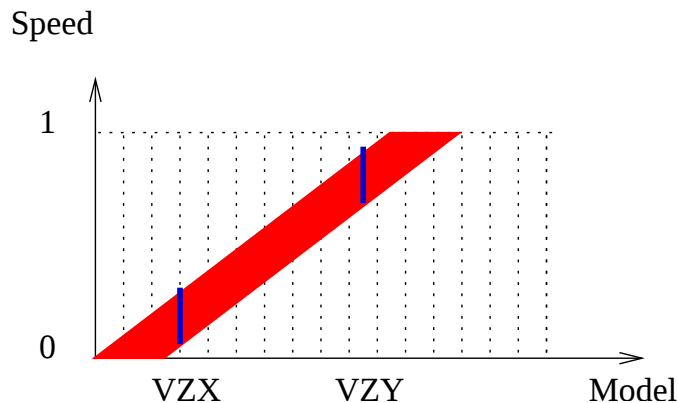
Truth value (Fuzziness level)

The value of youth of a 42 years-old man

- $V = 0$
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- $V = 0.5$
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- $V \in [0.2, 0.6]$
($0.2 \leq V \wedge V \leq 0.6$)
- $V \in [0.2, 0.5] \cup [0.8, 1]$
($0.2 \leq V \wedge V \leq 0.5$) $\vee$ ($0.8 \leq V \wedge V \leq 1$)

Fuzziness - Uncertainty

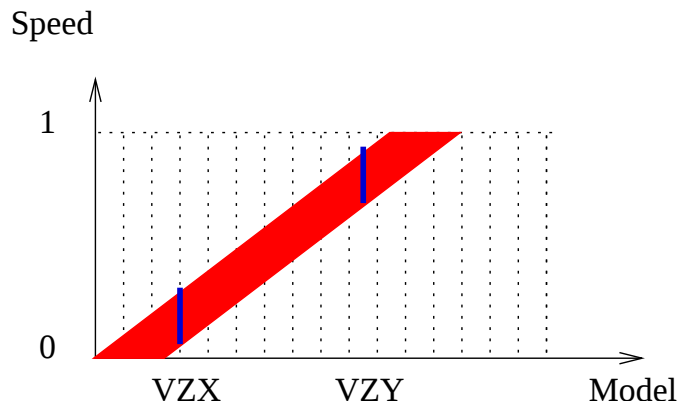
- New Laptop is a branch of computers with two laptop models (VZX and VZY). One model is very slow and the other one is very fast.



- VZX speed $[0.02, 0.08]$
- VZY speed $[0.75, 0.90]$

Fuzziness - Uncertainty

- New Laptop is a branch of computers with two laptop models (VZX and VZY). One model is very slow and the other one is very fast.



- VZX speed $[0.02, 0.08]$
- VZY speed $[0.75, 0.90]$

- If a client buys a New Laptop computer, the truth value, V , of its speed will be $[0.02, 0.08] \cup [0.75, 0.90]$
 $(0.02 \leq V \wedge V \leq 0.08) \vee (0.75 \leq V \wedge V \leq 0.90)$

Modeling Real World

- Knowledge:

Constraints ← { *Uncertainty*
Probability
Fuzziness
Incompleteness
Distributed knowledge

- Reasoning:

Prolog ← *Logic*

Fuzzy Prolog

- Existing Fuzzy Prolog systems:
 - Prolog-Elf
 - Fril Prolog
 - f-Prolog
- Our Fuzzy Prolog approach:
 - Truth Value (union of sub-intervals) $\mathcal{B}([0, 1])$
 - Aggregation operators (min, max, luka, ...)
 - CLP($\mathcal{R}$) based implementation

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Syntax

- If A is an atom, $A \leftarrow v$ is a *fuzzy fact*, where v , a truth value, is an element in $\mathcal{B}([0, 1])$ expressed as constraints over the domain $[0, 1]$.
- Let $A, B_1, \dots, B_n$ be atoms. A *fuzzy clause* is a clause of the form $A \leftarrow_F B_1 v_1, \dots, B_n v_n$ where F is an *aggregation operator* of truth values represented as constraints over the domain $[0, 1]$. The interval-aggregation induces a union-aggregation.
- A *fuzzy query* is a tuple $v \leftarrow A ?$ where A is an atom, and v is a variable (possibly instantiated) that represents a truth value in $\mathcal{B}([0, 1])$.

Aggregation Operators

- A function $f : [0, 1]^n \rightarrow [0, 1]$ that verifies $f(0, \dots, 0) = 0$, $f(1, \dots, 1) = 1$, and in addition it is monotonic and continuous, then it is called **aggregation operator**

Aggregation Operators

- A function $f : [0, 1]^n \rightarrow [0, 1]$ that verifies $f(0, \dots, 0) = 0$, $f(1, \dots, 1) = 1$, and in addition it is monotonic and continuous, then it is called **aggregation operator**
- Given an aggregation $f : [0, 1]^n \rightarrow [0, 1]$ an **interval-aggregation** $F : \mathcal{E}([0, 1])^n \rightarrow \mathcal{E}([0, 1])$ is defined as follows:

$$F([x_1^l, x_1^u], \dots, [x_n^l, x_n^u]) = [f(x_1^l, \dots, x_n^l), f(x_1^u, \dots, x_n^u)]$$

Union Aggregation

- Given an interval-aggregation $F : \mathcal{E}([0, 1])^n \rightarrow \mathcal{E}([0, 1])$ defined over intervals, a **union-aggregation** $\mathcal{F} : \mathcal{B}([0, 1])^n \rightarrow \mathcal{B}([0, 1])$ is defined over union of intervals as follows:

$$\mathcal{F}(B_1, \dots, B_n) = \cup \{F(\mathcal{E}_1, \dots, \mathcal{E}_n) \mid \mathcal{E}_i \in B_i\}$$

Interpretation

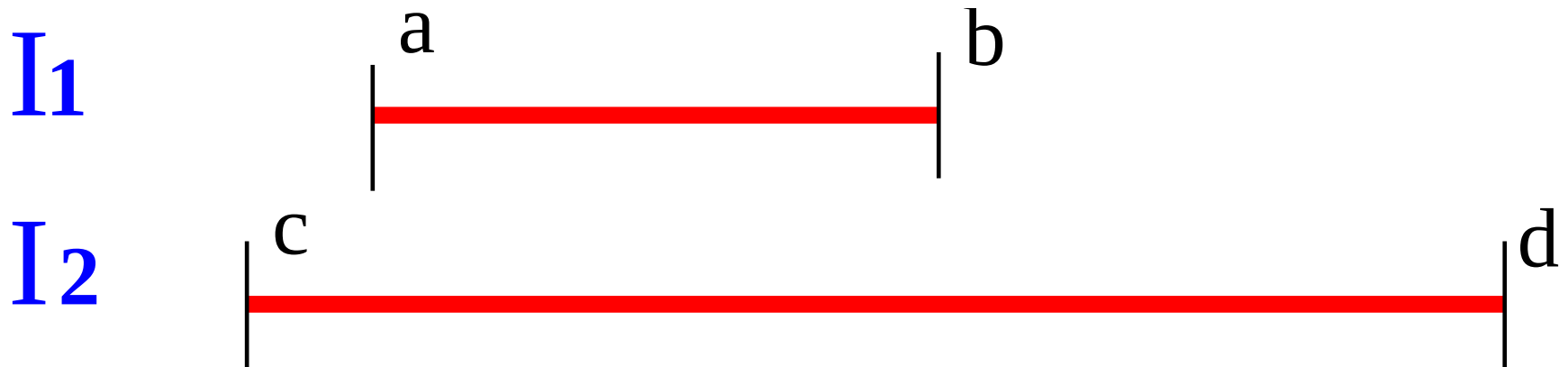
An *interpretation* I consists of the following:

1. a subset B_I of the *Herbrand Base*,
2. a mapping V_I , to assign a truth value, in $\mathcal{B}([0, 1])$, to each element of B_I .

The *Borel Algebra* $\mathcal{B}([0, 1])$ is a complete lattice under $\subseteq_{BI}$, that denotes *Borel inclusion*, and the *Herbrand Base* is a complete lattice under $\subseteq$, that denotes *set inclusion*, therefore a *set of all interpretations* forms a complete lattice under the relation $\sqsubseteq$ defined as follows.

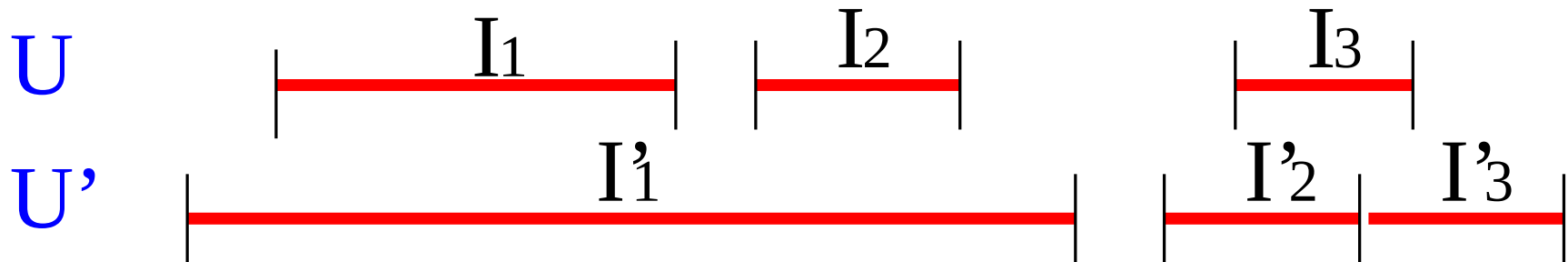
Interval Inclusion

Def:[interval inclusion $\subseteq_{II}$] Given two intervals $I_1 = [a, b]$, $I_2 = [c, d]$ in $\mathcal{E}([0, 1])$, $I_1 \subseteq_{II} I_2$ iff $c \leq a$ and $b \leq d$.



Borel Inclusion

Def:[Borel inclusion $\subseteq_{BI}$] Given two unions of intervals $U = I_1 \cup \dots \cup I_N$, $U' = I'_1 \cup \dots \cup I'_M$ in $\mathcal{B}([0, 1])$, $U \subseteq_{BI} U'$ if and only if $\forall x \in I_i, i \in 1..N$, $\exists I'_j \in U' . x \in I'_j$ where $j \in 1..M$.



Interpretation Inclusion - Valuation

Def:[interpretation inclusion $\sqsubseteq$] $I \sqsubseteq I'$ iff $B_I \subseteq B_{I'}$ and for all $B \in B_I$, $V_I(B) \subseteq_{B_I} V_{I'}(B)$, where $I = \langle B_I, V_I \rangle$, $I' = \langle B_{I'}, V_{I'} \rangle$ are interpretations.

Def:[valuation] A *valuation* σ of an atom A is an assignment of elements of U to variables of A . So $\sigma(A) \in B$ is a ground atom.

Model

Def: [model] Given an *interpretation* $I = \langle B_I, V_I \rangle$

- I is a *model* for a *fuzzy fact* $A \leftarrow v$, if for all valuation σ , $\sigma(A) \in B_I$ and $v \subseteq_{BI} V_I(\sigma(A))$.
- I is a *model* for a *clause* $A \leftarrow_F B_1, \dots, B_n$ when the following holds:
for all valuation σ , if $\sigma(B_i) \in B_I$, $1 \leq i \leq n$, and $v = \mathcal{F}(V_I(\sigma(B_1)), \dots, V_I(\sigma(B_n)))$ then $\sigma(A) \in B_I$ and $v \subseteq_{BI} V_I(\sigma(A))$, where $\mathcal{F}$ is the union aggregation obtained from F .
- I is a *model* of a *fuzzy program*, if it is a *model* for the facts and clauses of the program.

Semantics Equivalence

Given a program P , the three semantics:

1. **Least model** $lm(P)$, under the $\sqsubseteq$ ordering.
 2. **Declarative meaning** $lfp(T_P)$, least fixpoint for a consequence operator $T_P(I)$.
 3. **Success set** $SS(P)$ of a transitional system.
- are equivalent: $SS(P) = lfp(T_P) = lm(P)$.

Operational Semantics

- A sequence of transitions between different states of a system
- State: $\langle \textit{Goal}, \textit{Valuation}, \textit{Constraint} \rangle$
- Initial State: $\langle A, \emptyset, \textit{true} \rangle$
- Final State: $\langle \emptyset, \sigma, S \rangle$

Examples:

$\langle p(X, Y), \emptyset, \textit{true} \rangle, \dots, \langle \emptyset, \{X = 3, Y = 3\}, \textit{true} \rangle$

$\langle \textit{bachelor}(S, M), \emptyset, \textit{true} \rangle, \dots, \langle \emptyset, \{S = \textit{completed}\}, M \geq 5 \rangle$

Operational Semantics

A *transition* in the *transition system* is defined as:

1. $\langle A \cup a, \sigma, S \rangle \rightarrow \langle A\theta, \sigma \cdot \theta, S \wedge \mu_a = v \rangle$
if $h \leftarrow v$ is a fact of the program P , θ is the mgu of a and h , and μ_a is the truth variable for a , and $\text{solvable}(S \wedge \mu_a = v)$. ($\text{solvable}(c) \equiv c$ has solution in $[0, 1]$ of $\mathcal{R}$)
2. $\langle A \cup a, \sigma, S \rangle \rightarrow \langle (A \cup B)\theta, \sigma \cdot \theta, S \wedge c \rangle$
if $h \leftarrow_F B$ is a rule of the program P , θ is the mgu of a and h , c is the constraint that represents the truth value obtained applying the union-aggregator F on the truth variables of B , and $\text{solvable}(S \wedge c)$.
3. $\langle A \cup a, \sigma, S \rangle \rightarrow \text{fail}$ if none of the above are applicable.

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Fuzzy Programs Syntax

```
tall(john, V):~  
    [0.8, 0.9].  
tall(john, [0.8, 0.9]):~.
```

```
good_player(X, V):~min  
    tall(X, Vt),  
    swift(X, Vs).
```

```
f_digit :#  
    fuzzy digit/1.
```

```
not_small :#  
    fnot small/2.
```

Fuzzy $\rightarrow$ CLP($\mathcal{R}$) Translation

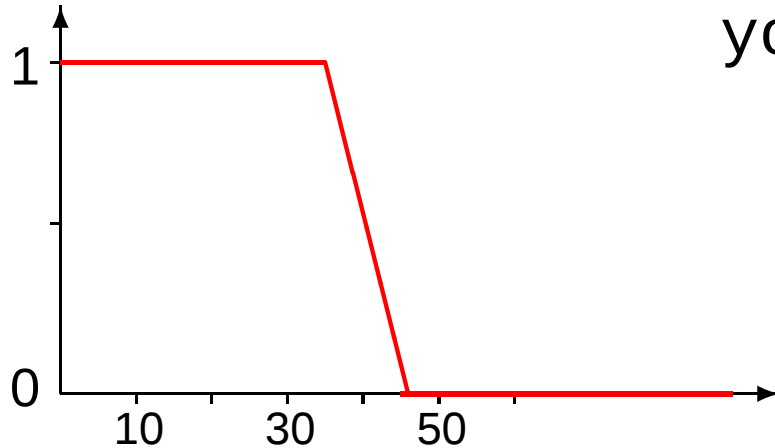
```
good_player(X,V):~min  
    tall(X,Vt),  
    swift(X,Vs).
```

```
good_player(X,V) :-  
    tall(X,Vt),  
    swift(X,Vs),  
    minim([Vt,Vs],V),  
    V.>=.0, V.<=.1.
```

```
not_small :#  
    fnot small/2.
```

```
not_small(X,V) :-  
    small(X,Vs),  
    V .=. 1 - Vs.
```

Syntactic Sugar



```
young :#  
fuzzy_predicate([ (0,1),  
                  (35,1),  
                  (45,0),  
                  (90,0) ]).
```

```
young(X,1):-  
    X .>=. 0,  
    X .<. 35.
```

```
young(X,V):-  
    X .>=. 35,  
    X .<. 45,  
    10*V.=.45-X.
```

```
young(X,0):-  
    X .>=. 45,  
    X .=<. 120.
```

Initial Evaluation

- Implementation over $\text{CLP}(\mathcal{R})$: SIMPLICITY
- Aggregation operator: GENERALITY
- Definition of new operators: FLEXIBILITY
- Using Prolog resolution: EFFICIENCY

Available implementation:

<http://clip.dia.fi.upm.es/Software/Ciao/>

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Combining Crisp and Fuzzy Logic

```
student(john).  
student(peter).
```

```
-----  
age_about_15(john,1):~.  
age_about_15(susan,0.7):~.  
age_about_15(nick,0):~.
```

```
-----  
teenager_student(X,V):~  
    student(X),           % CRISP  
    age_about_15(X,Va).% FUZZY
```

```
?- student(john).
```

```
yes
```

```
?- student(nick).
```

```
no                FALSE
```

```
?- age_about_15(john,V).
```

```
V = 1
```

```
?- age_about_15(nick,V).
```

```
V = 0
```

```
?- age_about_15(peter,V).
```

```
no                UNKNOWN
```

```
?- teenager_student(john,V).
```

```
V .=. 1
```

```
?- teenager_student(susan,V).
```

```
V .=. 0
```

```
?- teenager_student(peter,V).
```

```
no                UNKNOWN
```

Solution: Default Knowledge

```
student(john).
student(peter).

:-default(f_student/2,0).

f_student(X,1):-
    student(X).

-----

:-default(age_about_15/2,[0,1]).

age_about_15(john,1):~.
age_about_15(susan,0.7):~.
age_about_15(nick,0):~.

-----

:-default(teenager_student/2,[0,1]).

teenager_student(X,V):~
    f_student(X,Vs),
    age_about_15(X,Va).
```

```
?- f_student(john,V).
V = 1
?- f_student(nick,V).
V = 0                FALSE

?- age_about_15(john,V).
V = 1
?- age_about_15(nick,V).
V = 0
?- age_about_15(peter,V).
V .>=. 0, V .<=. 1    UNKNOWN

?- teenager_student(john,V).
V .=. 1
?- teenager_student(susan,V).
V .=. 0
?- teenager_student(peter,V).
V .>=. 0, V .<=. 1    UNKNOWN
```

Default Value

We assume there is a function *default* which implement the Default Knowledge Assumptions. It assigns an element of $\mathcal{B}([0, 1])$ to each element of the Herbrand Base.

- If the **Closed World Assumption** is used, then $default(A) = [0, 0]$ for all A in Herbrand Base.
- If **Open World Assumption** is used instead, $default(A) = [0, 1]$ for all A in Herbrand Base.

Interpretation

An *interpretation* I consists of the following:

1. a subset B_I of the *Herbrand Base*,
2. a mapping V_I , to assign
 - (a) a truth value, in $\mathcal{B}([0, 1])$, to each element of B_I , or
 - (b) $default(A)$, if A does not belong to B_I .

Operational Semantics

A *transition* in the *transition system* is defined as:

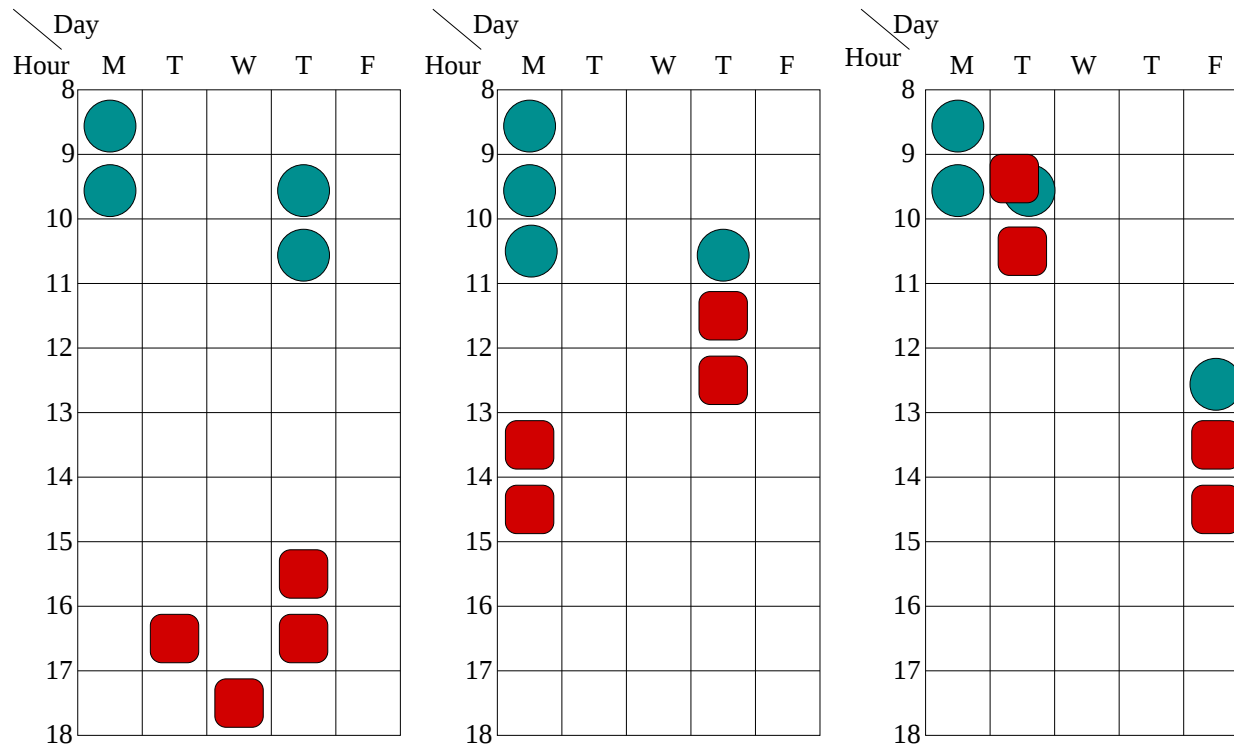
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if $h \leftarrow v$ is a fact of the program P , ...
2. $\langle A \cup a, \sigma, S \rangle \rightarrow \langle (A \cup B)\theta, \sigma \cdot \theta, S \wedge c \rangle$
if $h \leftarrow_F B$ is a rule of the program P , ...
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1. $\langle A \cup a, \sigma, S \rangle \rightarrow \langle A\theta, \sigma \cdot \theta, S \wedge \mu_a = v \rangle$
if $h \leftarrow v$ is a fact of the program P , ...
2. $\langle A \cup a, \sigma, S \rangle \rightarrow \langle (A \cup B)\theta, \sigma \cdot \theta, S \wedge c \rangle$
if $h \leftarrow_F B$ is a rule of the program P , ...
3. $\langle A \cup a, \sigma, S \rangle \rightarrow \text{fail}$ if none of the above are applicable.
 $\langle A \cup a, \sigma, S \rangle \rightarrow \langle A, \sigma, S \wedge \mu_a = v \rangle$
if none of the above are applicable and
 $\text{solvable}(S \wedge \mu_a = v)$ where $\mu_a = \text{default}(a)$.

Example (I) - Shifts Compatibility

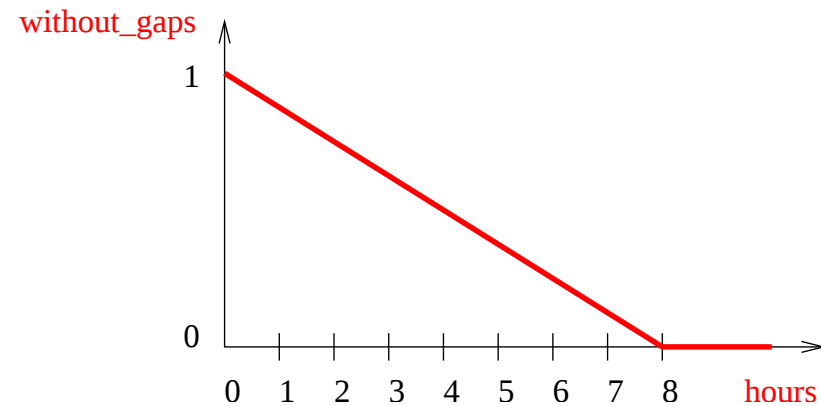
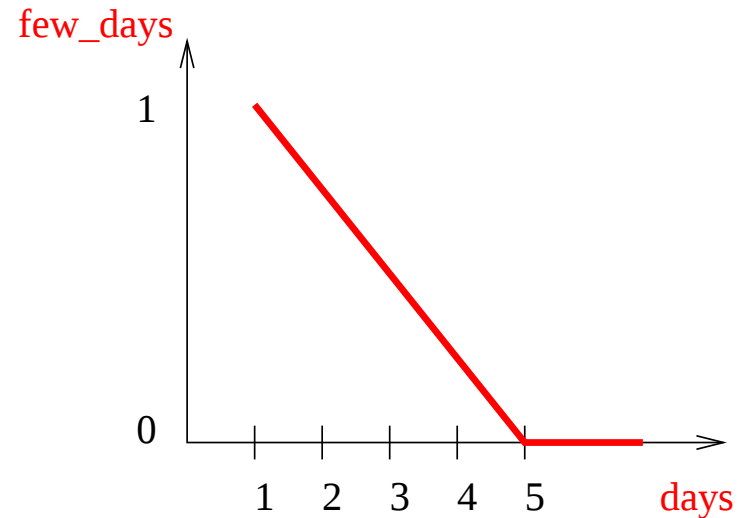


Timetables of compatible shifts

 Shift T1
 Shift T2

Example (II) - Crisp and Fuzzy

```
compatible(T1,T2,V):~ min
    correct_shift(T1),
    correct_shift(T2),
    disjoint(T1,T2),
    append(T1,T2,T),
    number_of_days(T,D),
    few_days(D,Vf),
    number_of_free_hours(T,H),
    without_gaps(H,Vw).
```



Example (III) - Default Values

```
f_correct_shift(T,1):- correct_shift(T).
```

```
:- default(f_correct_shift/2,[0,0]).      % CWA
```

```
f_disjoint(T1,T2,1):- disjoint(T1,T2).
```

```
:- default(f_disjoint/3,[0,0]).          % CWA
```

```
few_days(D,V):- ...
```

```
:- default(few_days/2,[0.25,0.75]).      % DEFAULT
```

```
without_gaps(H,V):- ...
```

```
:- default(without_gaps/2,[0,1]).         % OWA
```

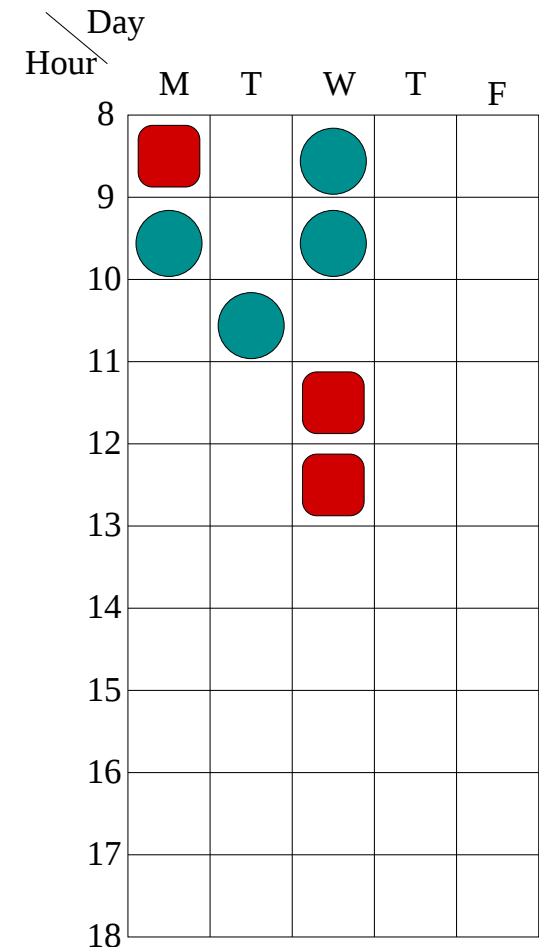
Example (IV) - Constructive Answers

?- compatible(
 [(mo, 9), (tu, 10), (we, 8), (we, 9)],
 [(mo, 8), (we, 11), (we, 12), (D, H)],
 V), V .>. 0.7 .

V = 0.9, D = we, H = 10 ? ;

V = 0.75, D = mo, H = 10 ? ;

no



Evaluation

- Representation of real problems: INCOMPLETENESS
- Crisp + Fuzzy logic: EXPRESSIVITY
- $[0, 1]$ to represent total uncertainty ($0 \leq v \wedge v \leq 1$). Lack of information do not stop the evaluation: ACCURACY
- Provides answers: CONSTRUCTIVE

Evaluation and Further Work

- Representation of real problems: INCOMPLETENESS
- Crisp + Fuzzy logic: EXPRESSIVITY
- $[0, 1]$ to represent total uncertainty ($0 \leq v \wedge v \leq 1$). Lack of information do not stop the evaluation: ACCURACY
- Provides answers: CONSTRUCTIVE

Further Work:

- Constructive negative queries
- Discrete fuzzy sets
- Applications (collaborative agents)

Overview

- Basics

- Introduction
- Description
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- Extensions

- Incompleteness
- **Constructive negative Queries**
- Discrete Fuzzy Sets
- Collaborative Fuzzy Agents
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Negation in Prolog - As Failure

student(john).
student(peter).

POSITIVE QUERIES

?- student(john).
yes
?- student(rick).
no

?- student(X).
X = john ? ;
X = peter ? ;
no

NEGATIVE QUERIES

?- +\ student(john).
no
?- +\ student(rick).
yes

?- +\ student(X).
no

Negation in Prolog - Constructive

student(john).
student(peter).

POSITIVE QUERIES

?- student(john).
yes
?- student(rick).
no

?- student(X).
X = john ? ;
X = peter ? ;
no

NEGATIVE QUERIES

?- +\ student(john).
no
?- +\ student(rick).
yes

?- neg(student(X)).
X = rick ?;
X = anne ?;
X = rose ?;
...

Negation in Prolog - Constructive

student(john).
student(peter).

POSITIVE QUERIES

?- student(john).
yes
?- student(rick).
no

?- student(X).
X = john ? ;
X = peter ? ;
no

NEGATIVE QUERIES

?- +\ student(john).
no
?- +\ student(rick).
yes

?- neg(student(X)).
X != john, X != peter ?;
no

Constructive Negation

?- neg(member(X, [1, 2])).

$X \neq 1, X \neq 2 ?;$

no

?- neg((X $\neq$ 0, member(X, [1, 2]))).

$X = 0 ?;$

$X \neq 0, X \neq 1, X \neq 2 ?;$

no

Fuzzified Crisp Predicates

```
student(john).  
student(peter).  
-----
```

```
f1_student(X,V):-  
    student(X),!,  
    V .=. 1.  
f1_student(X,0).  
-----
```

```
f2_student(X,V):-  
    student(X),  
    V .=. 1.  
f2_student(X,V):-  
    neg(student(X)),  
    V .=. 0.
```

```
?- f1_student(X,1).  
X = john ? ;  
X = peter ? ;  
no  
-----
```

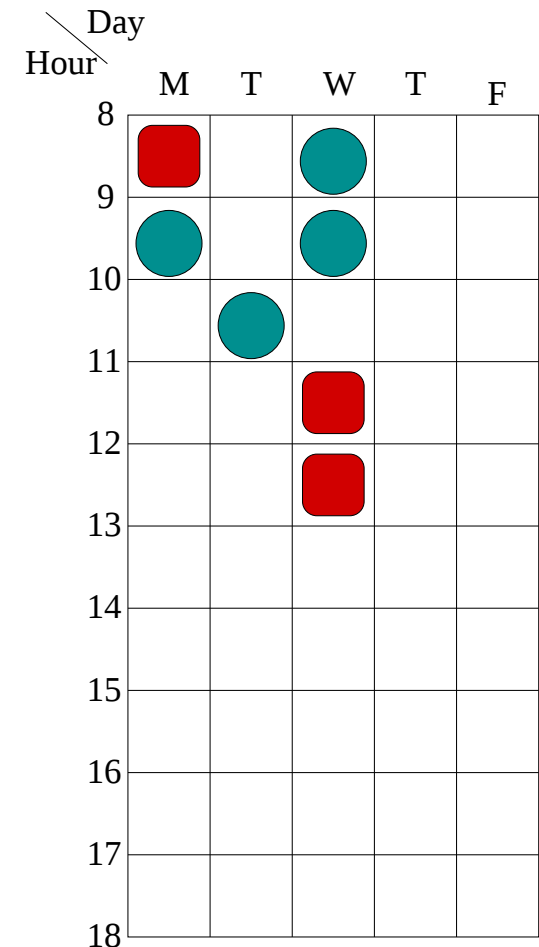
```
?- f1_student(X,0).  
no  
-----
```

```
?- f2_student(X,0).  
X != john, X != peter ? ;  
no
```

Example (V) - Constructive Answers

```
?- compatible(  
    [(mo, 9), (tu, 10),  
     (we, 8), (we, 9)],  
    [(mo, 8), (we, 11),  
     (we, 12), (D, 10)], V),  
    V .>. 0 .
```

```
D =/= tu ? ;  
no
```



Overview

- Basics

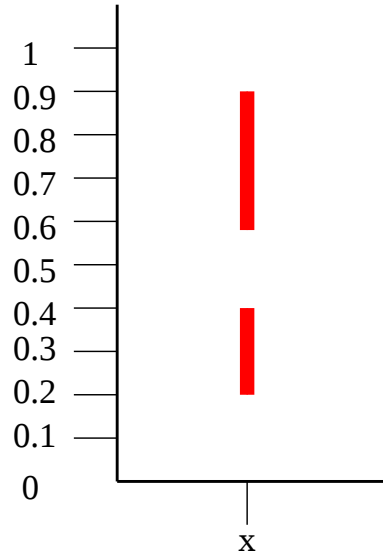
- Introduction
- Description
- Implementation

- Extensions

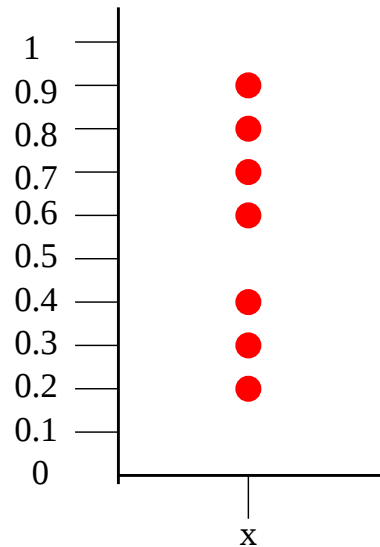
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Borel Alg. vs Discrete Borel Alg.



Borel Algebra



Discrete Borel Algebra

A *discrete-interval*

$[X_1, X_N]_d$ is a set of a finite number of values, $\{X_1, X_2, \dots, X_{N-1}, X_N\}$, between X_1 and X_N , $0 \leq X_1 \leq X_N \leq 1$, such that

$\exists 0 < \epsilon < 1. X_i = X_{i-1} + \epsilon, i \in \{2..N\}.$

Discrete Union Aggregation

- Given a discrete-interval-aggregation $F : \mathcal{E}_d([0, 1])^n \rightarrow \mathcal{E}_d([0, 1])$ defined over discrete-intervals, a **discrete-union-aggregation** $\mathcal{F} : \mathcal{B}_d([0, 1])^n \rightarrow \mathcal{B}_d([0, 1])$ is defined over union of discrete-intervals as follows:

$$\mathcal{F}(B_1, \dots, B_n) = \cup \{F(\mathcal{E}_{d,1}, \dots, \mathcal{E}_{d,n}) \mid \mathcal{E}_{d,i} \in B_i\}$$

Introduction to CLP($\mathcal{FD}$)

- CLP($\mathcal{FD}$): (Arithmetical) constraints over Finite Domains $\mathcal{FD}$: Each variable ranges over a finite set of integers
- *Resolution* is a combination of:
 - **Propagation**: excludes problem inconsistent values from the range of the variables (deterministic)
 - **Labeling**: assigns values to variables (expensive search process which fires more propagation)

Introduction to CLP($\mathcal{FD}$)

- Example:

```
main(X,Y,Z) :-  
    [X,Y,Z] in 1..5,  
    X-Y .=. 2*Z,  
    X+Y .≥. Z,  
    labeling([X,Y,Z]).
```

Fuzzy $\rightarrow$ CLP($\mathcal{FD}$) Translation

```
youth(45,V) :: ~  
  [0.2,0.5] v [0.8,1]
```

```
good_player(X,V)::~min  
  tall(X,Vt),  
  swift(X,Vs).
```

```
youth(45,V) :-  
  V in 2..5,  
  V in 8..10.
```

```
good_player(X,V) :-  
  tall(X,Vt),  
  swift(X,Vs),  
  minim([Vt,Vs],V),  
  V in 0..100.
```

Overview

- Basics

- Introduction
- Description
- Implementation

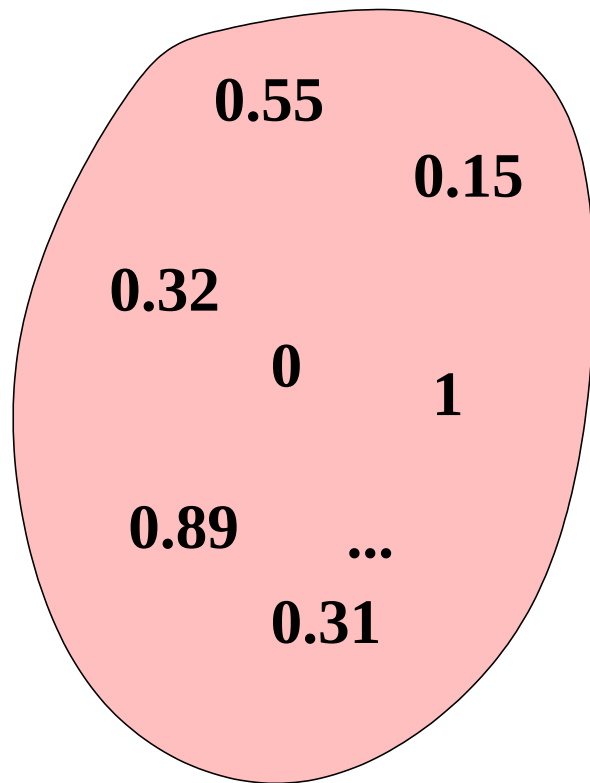
- Extensions

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Constraint Satisfaction Problems (CSP)

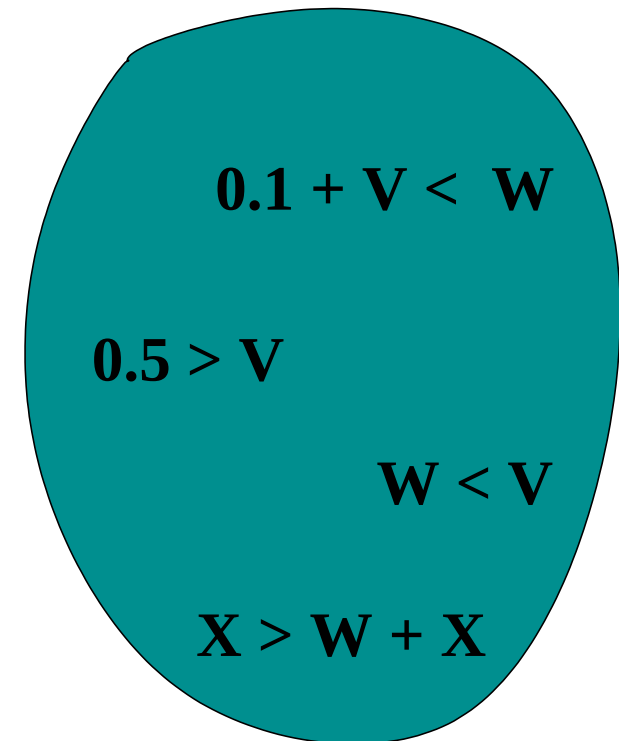
CANDIDATES



GOAL

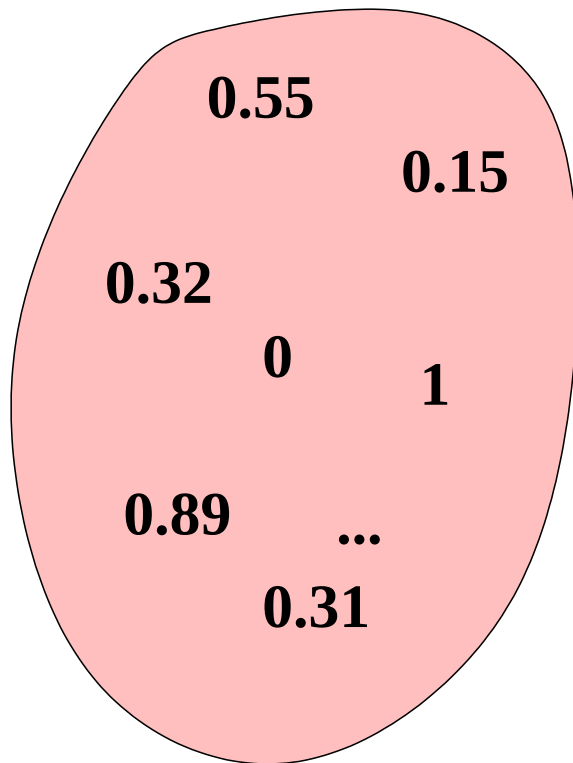
$V = ?$

CONSTRAINTS



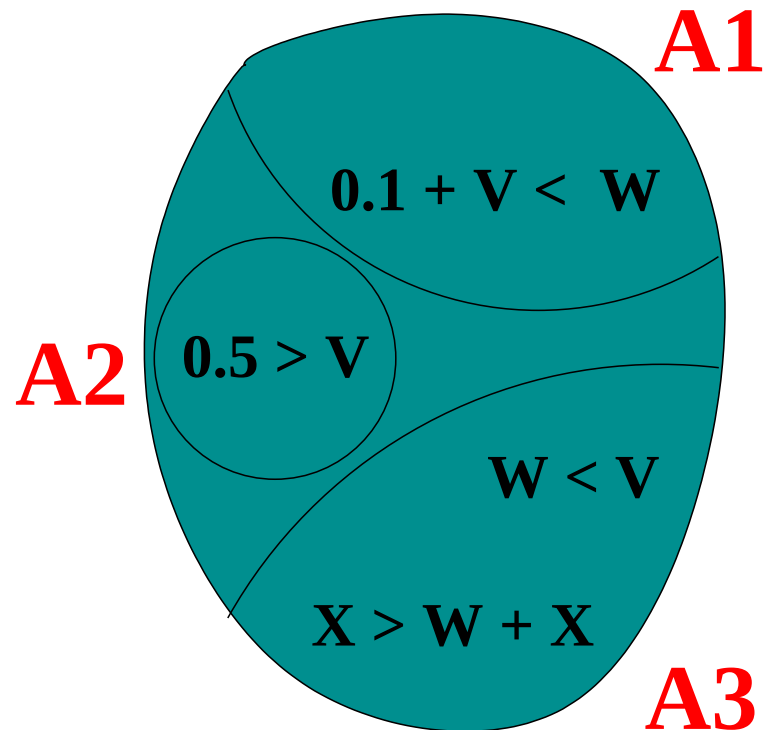
Distributed CSP

CANDIDATES



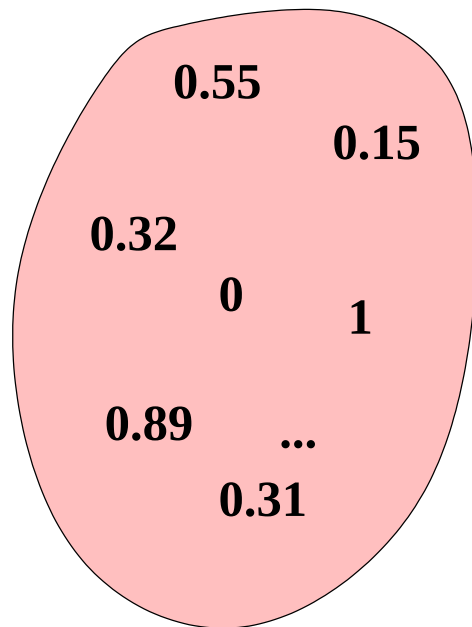
GOAL
 $V = ?$

CONSTRAINTS



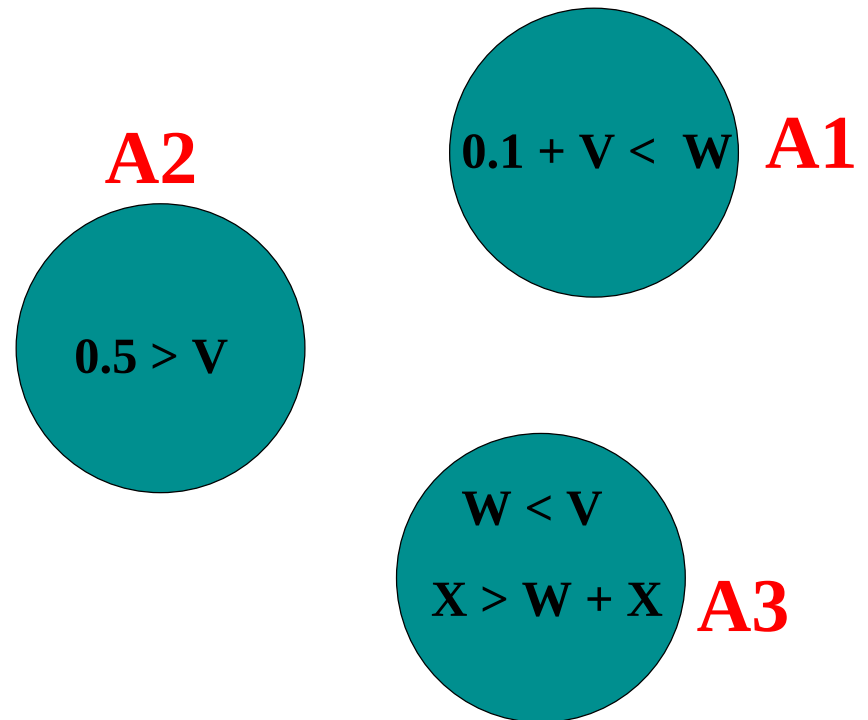
Distributed CSP (I)

CANDIDATES



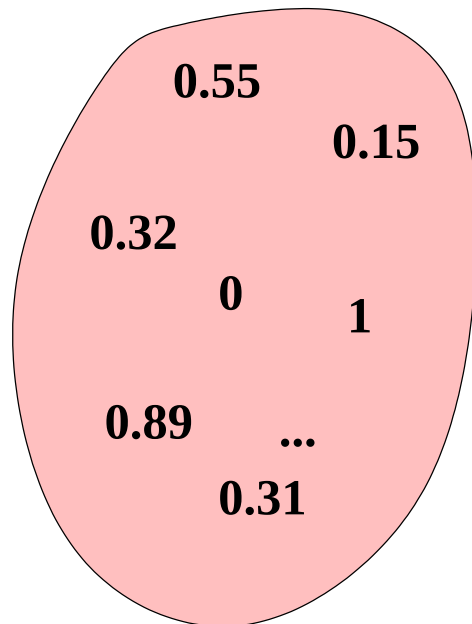
GOAL
 $V = ?$

CONSTRAINTS – AGENTS

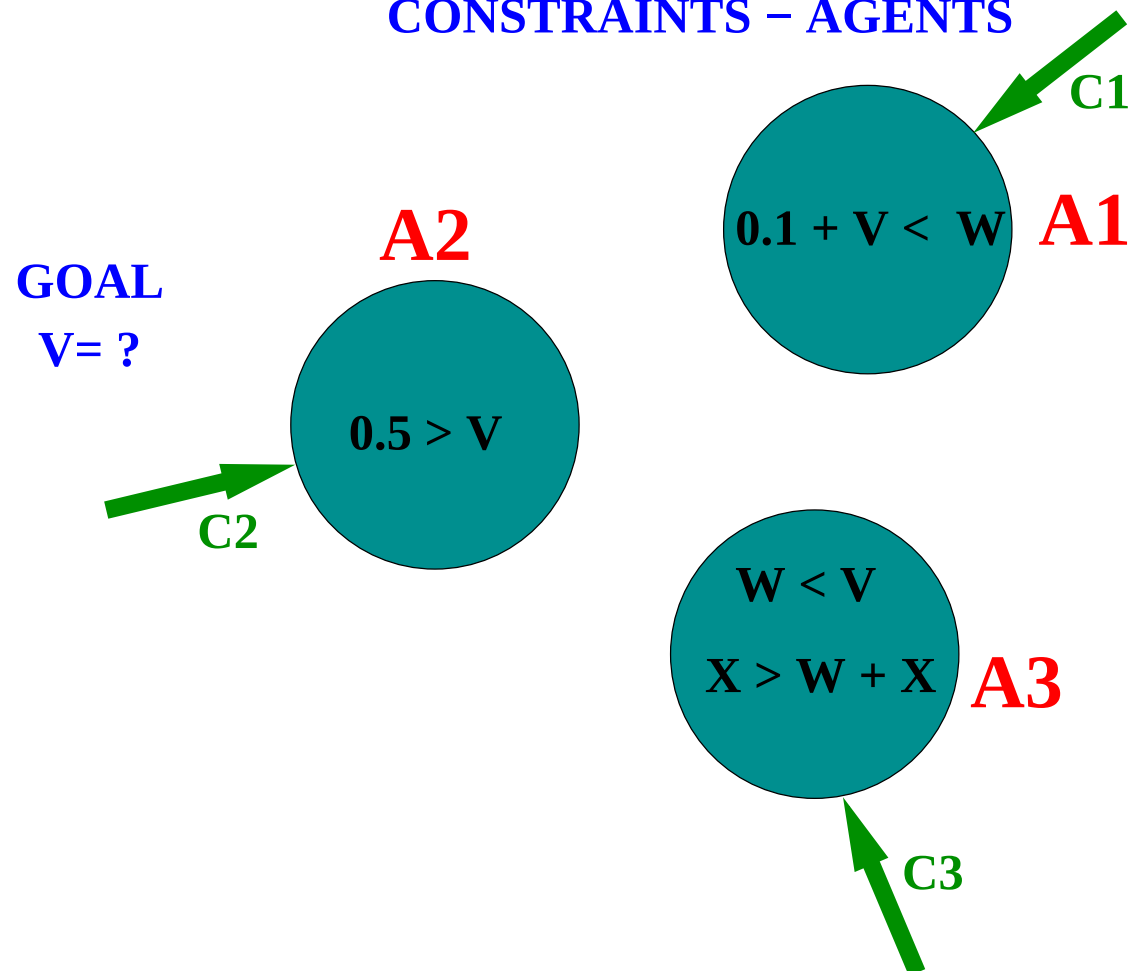


Distributed CSP (II)

CANDIDATES

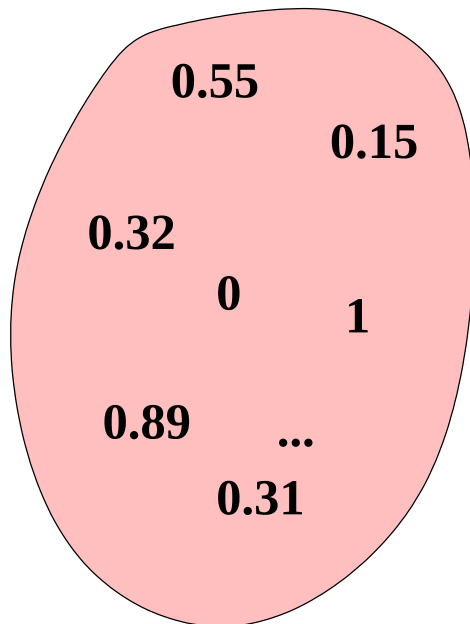


CONSTRAINTS – AGENTS

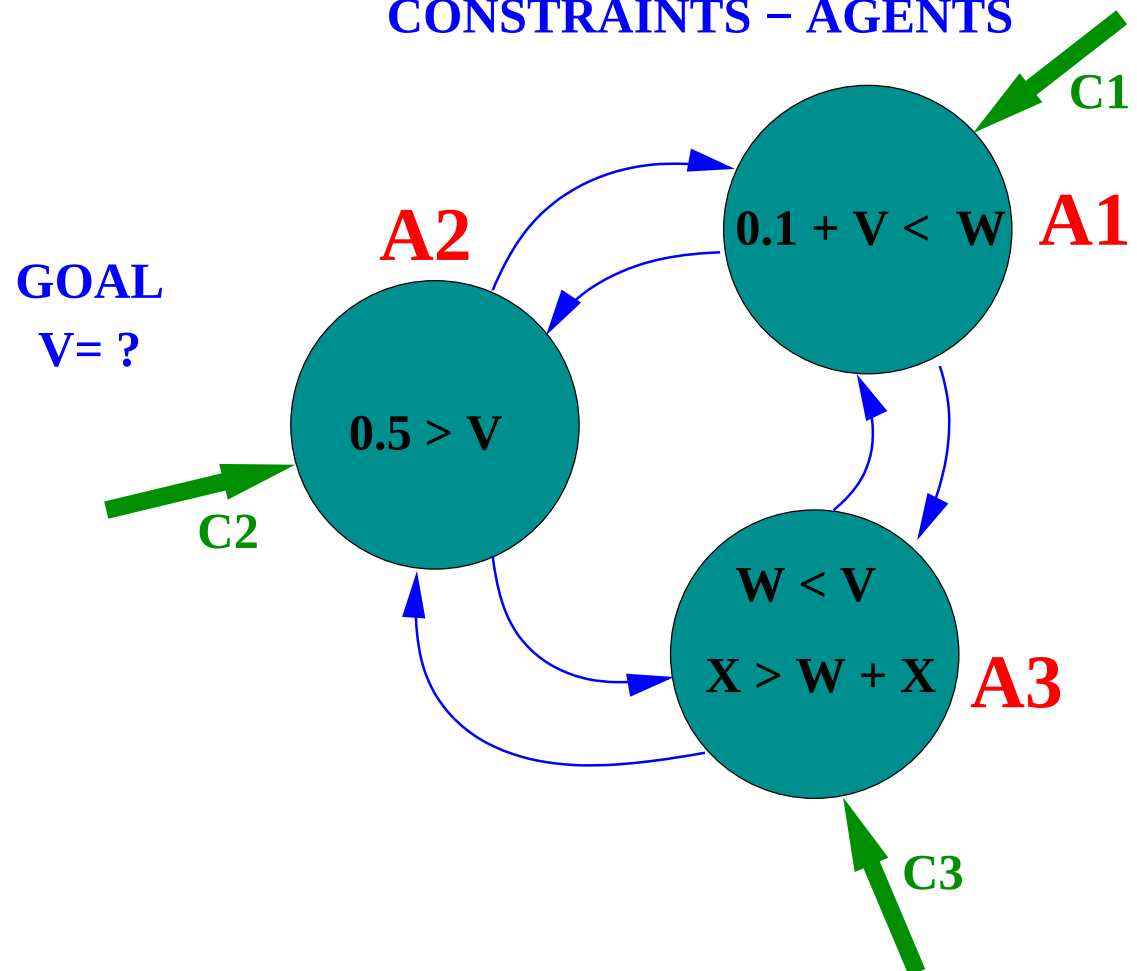


Distributed CSP (III)

CANDIDATES



CONSTRAINTS – AGENTS



Asynchronous Backtracking Algorithm

- Each agent owns exactly one variable and asynchronously assigns a value to its variable and sends it to the other agents (for evaluation)
- Each agent has a partial knowledge of the problem determined by the agents connected to it: *agent view*
- Messages exchanged:
 - *ok?* : assignment made by the agent
 - *nogood* : *agent view* which detects an inconsistency
 - *ack* : acknowledgement ($\equiv$ consistency)

Extended ABT (I)

- CLP($\mathcal{FD}$) resolution simplifies having **multiple variables** in each agent
- Coordination between distributed propagation and labeling
- We use the **Chandy-Lamport** algorithm for detecting when an agent is stable if:
 - It has no queued message
 - Its *agent view* is complete (there are no messages in transit)

Extended ABT (II)

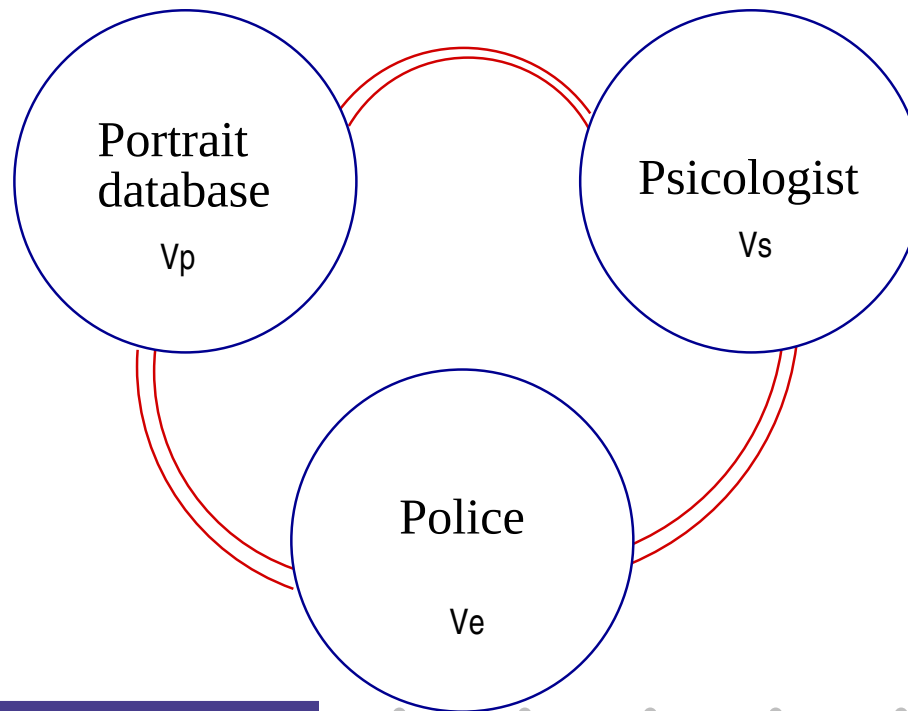
- Distributed propagation ends when termination is detected (all agents are in a stable state) → obtains a **global fixpoint** without inconsistent values
- We use **Dijkstra-Scholten** algorithm that provide a **reduction** of the **search** space → minimal exchange of propagation messages: minimal spanning tree

Collaborative Fuzzy Problems

- Collaborative Fuzzy Problems can be modelled using a combination of:
 - Discrete Fuzzy Prolog
 - An implementation of Extended ABT (for distributed reasoning)
- $CLP(\mathcal{FD})$ is the link between these components
- This work has been implemented in Ciao Prolog

Example (I)

- Criminal identification of suspects
- Distributed knowledge about:
 - physical aspects (V_p)
 - psychical aspects (V_s)
 - evidences (V_e)



Example (II)

- Discrete **fuzzy** program:

```
suspect(Person, V) ::~ inter_m  
    allocate_vars([Vp, Vs, Ve]),  
    physically_suspect(Person, Vp, Vs),  
    psychically_suspect(Person, Vs, Vp),  
    evidences(Person, Ve, Vp, Vs).
```

- Transformed **CLP($\mathcal{FD}$)** program:

```
suspect(Person, V) :-  
    allocate_vars([Vp, Vs, Ve]),  
    V in 0..10,  
    physically_suspect(Person, Vp, Vs),  
    psychically_suspect(Person, Vs, Vp),  
    evidences(Person, Ve, Vp, Vs),  
    inter_m([Vp, Vs, Ve], V).
```

Distributed Knowledge

- **Partial knowledge** stored in each agent is formulated in terms of **constraint expressions**

```
Cp:   physically_suspect(Person, Vp, Vs) :-  
        scan_portrait_database(Person, Vp),  
        Vp * Vs .>=. 50 @ a1.
```

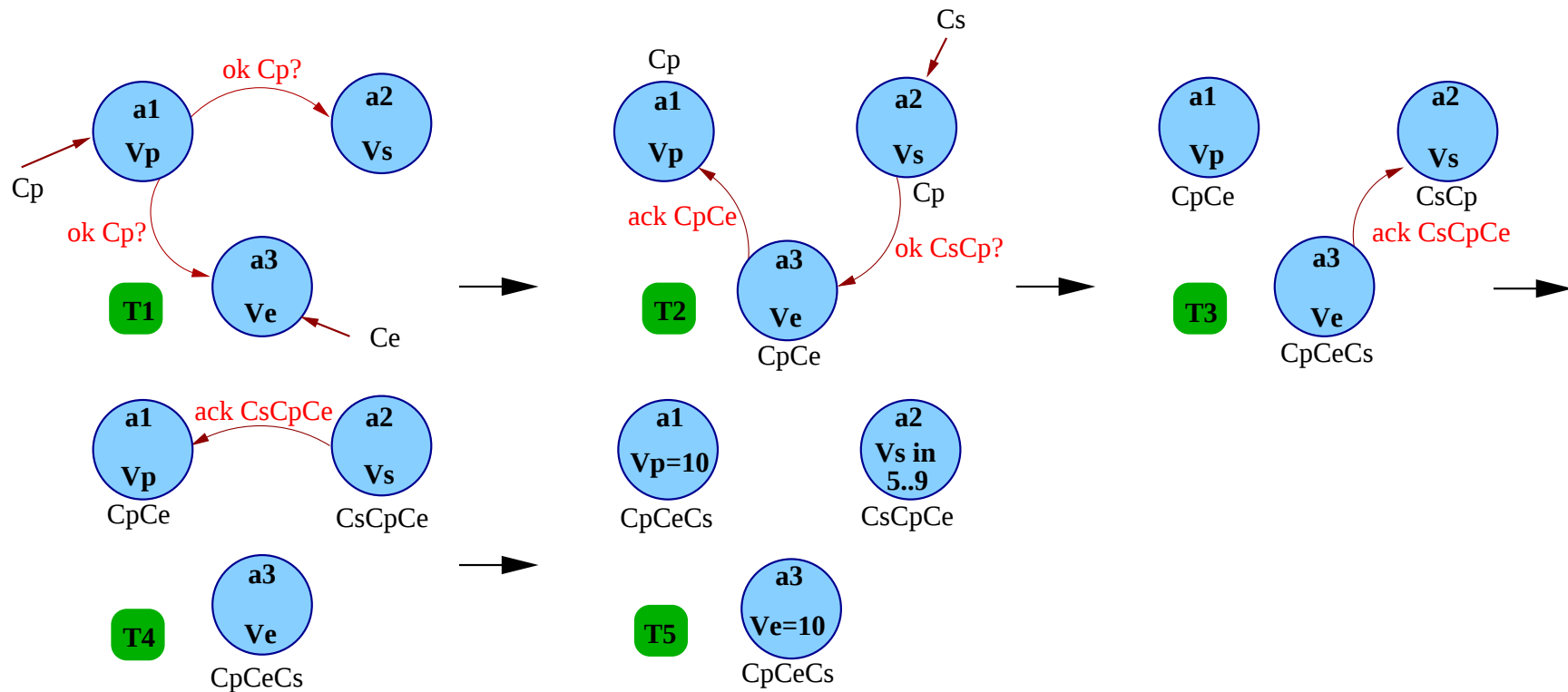
```
Cs:   psychically_suspect(Person, Vs, Vp) :-  
        psicologist_diagnostic(Person, Vs),  
        Vs .<. Vp @ a2.
```

```
Ce:   evidences(Person, Ve, Vp, Vs) :-  
        police_database(Person, Ve),  
        (Ve .>=. Vp,  
         Ve .>=. Vs) @ a3.
```

```
scan_portrait_database(peter, Vp) :- Vp in 4..10.  
psicologist_diagnostic(peter, Vs) :- Vs in 3..10.  
police_database(peter, Ve) :- Ve in 7..10.
```

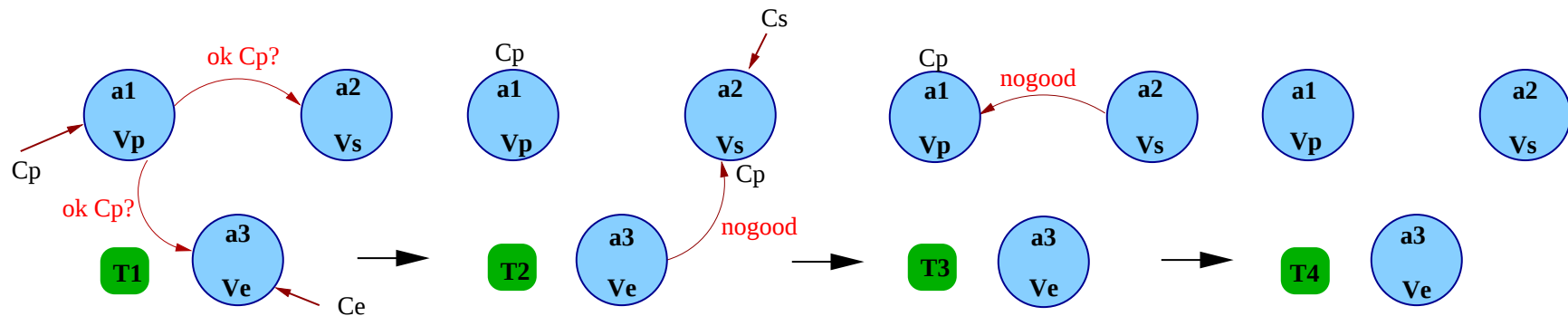
Agent Interaction (I)

- Collaborative fuzzy agents interaction for suspect(peter, V):



Agent Interaction (II)

- Collaborative fuzzy agents interaction for suspect(jane, V):



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Multi-adjoint logic

- Rules with a truth degree of **credibility**

$$\langle \mathcal{R}; \alpha \rangle$$

- Fuzzy Prolog rule syntax

$fpred(args) \text{ cred } (aggrC, \alpha) : \sim aggrO$

$fpred1(args1), \dots, fpredN(argsN).$

- Example

$good_player(J) \text{ cred } (prod, 0.8) : \sim prod$
 $swift(J), tall(J), experience(J).$

Multi-adjoint logic

- Fuzzy Prolog fact syntax

fpred(args) value truth_value.

- Example

```
experience(john) value 0.9 .  
experience(karl) value 0.9 .  
experience(mike) value 0.9 .  
experience(lebron) value 0.4 .  
experience(deron) value 0.3 .
```

Default values

- Represent incomplete information
- Fuzzy Prolog default value syntax

: – `default (fpred/arity, default_value)`.

- Example

```
:- default (experience/1, 0.9).  
experience(lebron) value 0.4 .  
experience(deron) value 0.3 .
```

Type Properties

- Fuzzy Prolog type properties syntax

$: - \text{prop } \textit{type_name}/\textit{arity}.$

$: - \text{set_prop } \textit{fpred}(\textit{args}) \Rightarrow \textit{type_name}(\textit{args}).$

- Example

```
:- prop typePlayer/1.
```

```
typePlayer(john).
```

```
...
```

```
typePlayer(deron).
```

```
:- set_prop experience(J) => typePlayer(J).
```

Complete Example I

```
:- module(good_player,_, [rfuzzy]).  
:- prop typePlayer/1.
```

```
(good_player(J) cred (prop,0.8)) :~ prop  
    swift(J), tall(J), experience(J).
```

```
typePlayer(john).
```

```
...
```

```
typePlayer(deron).
```

```
:- set_prop experience(J) >= typePlayer(J).
```

Complete Example II

```
:- default (experience/1, 0.9).  
experience(lebron) value 0.4 .  
experience(deron) value 0.3 .
```

```
:- default (tall/1, 0.6).  
tall(john) value 0.4 .  
tall(karl) value 0.8 .
```

```
:- default (swift/1, 0.7).  
swift(john) value 1 .
```

Fuzzy Queries

- Fuzzy Prolog queries syntax

$? - \text{fpred}(\text{args}, V).$

- Example

`?- good_player(john, V).`

`V= 0.288 ?;`

`no`

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Work proposals

- “Models of Inexact Reasoning” work
 - Modeling a Problem with Fuzzy/RFuzzy Prolog
 - Semantics for Fuzzy queries language
 - ...
- Practical Project / Master thesis
 - Implementation of credibility with intervals
 - Fuzzy web interface
 - New Fuzzy Prolog version with credibility
 - Comparison with other Fuzzy Prolog (tool & semantics)
 - Credibility with intervals
 - ...

Work “Models of Inexact Reasoning”

- Choose a topic (with the acknowledge of the professor)
- Send **Feb 20th** by e-mail to [*susana@fi.upm.es*](mailto:susana@fi.upm.es)
 - Source files of the report (better .tex, otherwise .doc)
 - Report with tests, explanations, etc. (.pdf)
 - Source files of the developed code (.pl)
 - Examples files (.pl)

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Fuzzy Prolog

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