

All Rewrite Rules

From Event-B

CAUTION! Any modification to this page shall be announced on the User mailing list!

This page groups together all the rewrite rules implemented (or planned for implementation) in the Rodin prover. The rules themselves can be found in their respective location (for editing purposes):

Conventions used in these tables are described in [The_Proving_Perspective_\(Rodin_User_Manual\)#Rewrite_rules](#)

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Set Rewrite Rules

Rules that are marked with a * in the first column are implemented in the latest version of Rodin. Rules without a * are planned to be implemented in future versions. Other conventions used in these tables are described in [The_Proving_Perspective_\(Rodin_User_Manual\)#Rewrite_Rules](#).

	Name	Rule	Side Cor
*	SIMP_SPECIAL_AND_BTRUE	$P \wedge \dots \wedge \top \wedge \dots \wedge Q \cong P \wedge \dots \wedge Q$	
*	SIMP_SPECIAL_AND_BFALSE	$P \wedge \dots \wedge \perp \wedge \dots \wedge Q \cong \perp$	
*	SIMP_MULTI_AND	$P \wedge \dots \wedge Q \wedge \dots \wedge Q \wedge \dots \wedge R \cong P \wedge \dots \wedge Q \wedge \dots \wedge R$	
*	SIMP_MULTI_AND_NOT	$P \wedge \dots \wedge Q \wedge \dots \wedge \neg Q \wedge \dots \wedge R \cong \perp$	
*	SIMP_SPECIAL_OR_BTRUE	$P \vee \dots \vee \top \vee \dots \vee Q \cong \top$	
*	SIMP_SPECIAL_OR_BFALSE	$P \vee \dots \vee \perp \vee \dots \vee Q \cong P \vee \dots \vee Q$	
*	SIMP_MULTI_OR	$P \vee \dots \vee Q \vee \dots \vee Q \vee \dots \vee R \cong P \vee \dots \vee Q \vee \dots \vee R$	
*	SIMP_MULTI_OR_NOT	$P \vee \dots \vee Q \vee \dots \vee \neg Q \wedge \dots \wedge R \cong \top$	
*	SIMP_SPECIAL_IMP_BTRUE_R	$P \Rightarrow \top \cong \top$	
*	SIMP_SPECIAL_IMP_BTRUE_L	$\top \Rightarrow P \cong P$	
*	SIMP_SPECIAL_IMP_BFALSE_R	$P \Rightarrow \perp \cong \neg P$	
*	SIMP_SPECIAL_IMP_BFALSE_L	$\perp \Rightarrow P \cong \top$	
*	SIMP_MULTI_IMP	$P \Rightarrow P \cong \top$	
*	SIMP_MULTI_IMP_NOT_L	$\neg P \Rightarrow P \cong P$	
*	SIMP_MULTI_IMP_NOT_R	$P \Rightarrow \neg P \cong \neg P$	
*	SIMP_MULTI_IMP_AND	$P \wedge \dots \wedge Q \wedge \dots \wedge R \Rightarrow Q \cong \top$	
*	SIMP_MULTI_IMP_AND_NOT_R	$P \wedge \dots \wedge Q \wedge \dots \wedge R \Rightarrow \neg Q \cong \neg(P \wedge \dots \wedge Q \wedge \dots \wedge R)$	
*	SIMP_MULTI_IMP_AND_NOT_L	$P \wedge \dots \wedge \neg Q \wedge \dots \wedge R \Rightarrow Q \cong \neg(P \wedge \dots \wedge \neg Q \wedge \dots \wedge R)$	
*	SIMP_MULTI_EQV	$P \Leftrightarrow P \cong \top$	
*	SIMP_MULTI_EQV_NOT	$P \Leftrightarrow \neg P \cong \perp$	
*	SIMP_SPECIAL_NOT_BTRUE	$\neg \top \cong \perp$	
*	SIMP_SPECIAL_NOT_BFALSE	$\neg \perp \cong \top$	
*	SIMP_NOT_NOT	$\neg \neg P \cong P$	
*	SIMP_NOTEQUAL	$E \neq F \cong \neg E = F$	
*	SIMP_NOTIN	$E \notin F \cong \neg E \in F$	

*	SIMP_NOTSUBSET	$E \not\subset F \hat{=} \neg E \subset F$	
*	SIMP_NOTSUBSETEQ	$E \not\subseteq F \hat{=} \neg E \subseteq F$	
*	SIMP_NOT_LE	$\neg a \leq b \hat{=} a > b$	
*	SIMP_NOT_GE	$\neg a \geq b \hat{=} a < b$	
*	SIMP_NOT_LT	$\neg a < b \hat{=} a \geq b$	
*	SIMP_NOT_GT	$\neg a > b \hat{=} a \leq b$	
*	SIMP_SPECIAL_NOT_EQUAL_FALSE_R	$\neg(E = \text{FALSE}) \hat{=} (E = \text{TRUE})$	
*	SIMP_SPECIAL_NOT_EQUAL_FALSE_L	$\neg(\text{FALSE} = E) \hat{=} (\text{TRUE} = E)$	
*	SIMP_SPECIAL_NOT_EQUAL_TRUE_R	$\neg(E = \text{TRUE}) \hat{=} (E = \text{FALSE})$	
*	SIMP_SPECIAL_NOT_EQUAL_TRUE_L	$\neg(\text{TRUE} = E) \hat{=} (\text{FALSE} = E)$	
*	SIMP_FORALL_AND	$\forall x.P \wedge Q \hat{=} (\forall x.P) \wedge (\forall x.Q)$	
*	SIMP_EXISTS_OR	$\exists x.P \vee Q \hat{=} (\exists x.P) \vee (\exists x.Q)$	
*	SIMP_EXISTS_IMP	$\exists x.P \Rightarrow Q \hat{=} (\forall x.P) \Rightarrow (\exists x.Q)$	
*	SIMP_FORALL	$\forall \dots, z, \dots.P(z) \hat{=} \forall z.P(z)$	Quantified identifiers than z do r in P
*	SIMP_EXISTS	$\exists \dots, z, \dots.P(z) \hat{=} \exists z.P(z)$	Quantified identifiers than z do r in P
*	SIMP_MULTI_EQUAL	$E = E \hat{=} \top$	
*	SIMP_MULTI_NOTEQUAL	$E \neq E \hat{=} \perp$	
*	SIMP_EQUAL_MAPSTO	$E \mapsto F = G \mapsto H \hat{=} E = G \wedge F = H$	
*	SIMP_EQUAL_SING	$\{E\} = \{F\} \hat{=} E = F$	
*	SIMP_SPECIAL_EQUAL_TRUE	$\text{TRUE} = \text{FALSE} \hat{=} \perp$	
*	SIMP_TYPE_SUBSETEQ	$S \subseteq Ty \hat{=} \top$	where Ty is expression
*	SIMP_SUBSETEQ_SING	$\{E\} \subseteq S \hat{=} E \in S$	where E single expr
*	SIMP_SPECIAL_SUBSETEQ	$\emptyset \subseteq S \hat{=} \top$	
*	SIMP_MULTI_SUBSETEQ	$S \subseteq S \hat{=} \top$	
*	SIMP_SUBSETEQ_BUNION	$S \subseteq A \cup \dots \cup S \cup \dots \cup B \hat{=} \top$	
*	SIMP_SUBSETEQ_BINTER	$A \cap \dots \cap S \cap \dots \cap B \subseteq S \hat{=} \top$	
*	DERIV_SUBSETEQ_BUNION	$A \cup \dots \cup B \subseteq S \hat{=} A \subseteq S \wedge \dots \wedge B \subseteq S$	
*	DERIV_SUBSETEQ_BINTER	$S \subseteq A \cap \dots \cap B \hat{=} S \subseteq A \wedge \dots \wedge S \subseteq B$	
*	SIMP_SPECIAL_IN	$E \in \emptyset \hat{=} \perp$	
*	SIMP_MULTI_IN	$B \in \{A, \dots, B, \dots, C\} \hat{=} \top$	
*	SIMP_IN_SING	$E \in \{F\} \hat{=} E = F$	
*	SIMP_MULTI_SETENUM	$\{A, \dots, B, \dots, B, \dots, C\} \hat{=} \{A, \dots, B, \dots, C\}$	
*	SIMP_SPECIAL_BINTER	$S \cap \dots \cap \emptyset \cap \dots \cap T \hat{=} \emptyset$	
*	SIMP_TYPE_BINTER	$S \cap \dots \cap Ty \cap \dots \cap T \hat{=} S \cap \dots \cap T$	where Ty is expression
*	SIMP_MULTI_BINTER	$S \cap \dots \cap T \cap \dots \cap T \cap \dots \cap U \hat{=} S \cap \dots \cap T \cap \dots \cap U$	
*	SIMP_MULTI_EQUAL_BINTER	$S \cap \dots \cap T \cap \dots \cap U = T \hat{=} T \subseteq S \cap \dots \cap U$	
*	SIMP_SPECIAL_BUNION	$S \cup \dots \cup \emptyset \cup \dots \cup T \hat{=} S \cup \dots \cup T$	
*	SIMP_TYPE_BUNION	$S \cup \dots \cup Ty \cup \dots \cup T \hat{=} Ty$	where Ty is expression
*	SIMP_MULTI_BUNION	$S \cup \dots \cup T \cup \dots \cup T \cup \dots \cup U \hat{=} S \cup \dots \cup T \cup \dots \cup U$	

*	SIMP_MULTI_EQUAL_UNION	$S \cup \dots \cup T \cup \dots \cup U = T \hat{=} S \cup \dots \cup U \subseteq T$	
*	SIMP_MULTI_SETMINUS	$S \setminus S \hat{=} \emptyset$	
*	SIMP_SPECIAL_SETMINUS_R	$S \setminus \emptyset \hat{=} S$	
*	SIMP_SPECIAL_SETMINUS_L	$\emptyset \setminus S \hat{=} \emptyset$	
*	SIMP_TYPE_SETMINUS	$S \setminus Ty \hat{=} \emptyset$	where Ty is an expression
*	SIMP_TYPE_SETMINUS_SETMINUS	$Ty \setminus (Ty \setminus S) \hat{=} S$	where Ty is an expression
*	SIMP_KUNION_POW	$\text{union}(\mathbb{P}(S)) \hat{=} S$	
*	SIMP_KUNION_POW1	$\text{union}(\mathbb{P}_1(S)) \hat{=} S$	
*	SIMP_SPECIAL_KUNION	$\text{union}(\{\emptyset\}) \hat{=} \emptyset$	
*	SIMP_SPECIAL_QUNION	$\bigcup x.\perp \mid E \hat{=} \emptyset$	
*	SIMP_SPECIAL_KINTER	$\text{inter}(\{\emptyset\}) \hat{=} \emptyset$	
*	SIMP_KINTER_POW	$\text{inter}(\mathbb{P}(S)) \hat{=} \emptyset$	
*	SIMP_SPECIAL_POW	$\mathbb{P}(\emptyset) \hat{=} \{\emptyset\}$	
*	SIMP_SPECIAL_POW1	$\mathbb{P}_1(\emptyset) \hat{=} \emptyset$	
*	SIMP_SPECIAL_CPROD_R	$S \times \emptyset \hat{=} \emptyset$	
*	SIMP_SPECIAL_CPROD_L	$\emptyset \times S \hat{=} \emptyset$	
	SIMP_COMPSET_EQUAL	$\{x, y \cdot x = E(y) \wedge P(y) \mid F(x, y)\} \hat{=} \{y \cdot P(y) \mid F(E(y), y)\}$	where x is not in E and in P
*	SIMP_COMPSET_IN	$\{x \cdot x \in S \mid x\} \hat{=} S$	where x is not in S
*	SIMP_COMPSET_SUBSETEQ	$\{x \cdot x \subseteq S \mid x\} \hat{=} \mathbb{P}(S)$	where x is not in S
*	SIMP_SPECIAL_COMPSET_BFALSE	$\{x \cdot \perp \mid x\} \hat{=} \emptyset$	
*	SIMP_SPECIAL_COMPSET_BTRUE	$\{x \cdot \top \mid E\} \hat{=} Ty$	where the Ty is a maplet combinatoric locally-bound pairwise-disjoint bound identifier
*	SIMP_SUBSETEQ_COMPSET_L	$\{x \cdot P(x) \mid E(x)\} \subseteq S \hat{=} \forall y \cdot P(y) \Rightarrow E(y) \in S$	where y is not in S
*	SIMP_IN_COMPSET	$F \in \{x, y, \dots \cdot P(x, y, \dots) \mid E(x, y, \dots)\} \hat{=} \exists x, y, \dots \cdot P(x, y, \dots) \wedge E(x, y, \dots) = F$	where x, y, \dots are not free
*	SIMP_IN_COMPSET_ONEPOINT	$E \in \{x \cdot P(x) \mid x\} \hat{=} P(E)$	Equivalent general simplification followed by Point Rule application the last corollary predicate
	SIMP_SUBSETEQ_COMPSET_R	$S \subseteq \{x \cdot P(x) \mid x\} \hat{=} \forall y \cdot y \in S \Rightarrow P(y)$	where y is not in $S, \{x \cdot P(x) \mid x\}$
*	SIMP_SPECIAL_OVERL	$r \triangleleft \dots \triangleleft \emptyset \triangleleft \dots \triangleleft s \hat{=} r \triangleleft \dots \triangleleft s$	
*	SIMP_SPECIAL_KBOOL_BTRUE	$\text{bool}(\top) \hat{=} \text{TRUE}$	
*	SIMP_SPECIAL_KBOOL_BFALSE	$\text{bool}(\perp) \hat{=} \text{FALSE}$	
	DISTRIB_SUBSETEQ_UNION_SING	$S \cup \{F\} \subseteq T \hat{=} S \subseteq T \wedge F \in T$	where F is an expression
*	DEF_FINITE	$\text{finite}(S) \hat{=} \exists n, f \cdot f \in 1..n \mapsto S$	
*	SIMP_SPECIAL_FINITE	$\text{finite}(\emptyset) \hat{=} \top$	

* SIMP_FINITE_SETENUM	$\text{finite}(\{a, \dots, b\}) \hat{=} \top$	
* SIMP_FINITE_BUNION	$\text{finite}(S \cup T) \hat{=} \text{finite}(S) \wedge \text{finite}(T)$	
SIMP_FINITE_UNION	$\text{finite}(\text{union}(S)) \hat{=} \text{finite}(S) \wedge (\forall x \cdot x \in S \Rightarrow \text{finite}(x))$	
SIMP_FINITE_QUONION	$\text{finite}(\bigcup x \cdot P \mid E) \hat{=} \text{finite}(\{x \cdot P \mid E\}) \wedge (\forall x \cdot P \Rightarrow \text{finite}(E))$	
* SIMP_FINITE_POW	$\text{finite}(\mathbb{P}(S)) \hat{=} \text{finite}(S)$	
* DERIV_FINITE_CPROD	$\text{finite}(S \times T) \hat{=} S = \emptyset \vee T = \emptyset \vee (\text{finite}(S) \wedge \text{finite}(T))$	
* SIMP_FINITE_CONVERSE	$\text{finite}(r^{-1}) \hat{=} \text{finite}(r)$	
* SIMP_FINITE_UPTO	$\text{finite}(a..b) \hat{=} \top$	
* SIMP_FINITE_ID	$\text{finite}(\text{id}) \hat{=} \text{finite}(S)$	where id h $\mathbb{P}(S \times S)$
* SIMP_FINITE_ID_DOMRES	$\text{finite}(E \triangleleft \text{id}) \hat{=} \text{finite}(E)$	
* SIMP_FINITE_PRJ1	$\text{finite}(\text{prj}_1) \hat{=} \text{finite}(S \times T)$	where prj_1 type $\mathbb{P}(S \times T)$
* SIMP_FINITE_PRJ2	$\text{finite}(\text{prj}_2) \hat{=} \text{finite}(S \times T)$	where prj_2 type $\mathbb{P}(S \times T)$
* SIMP_FINITE_PRJ1_DOMRES	$\text{finite}(E \triangleleft \text{prj}_1) \hat{=} \text{finite}(E)$	
* SIMP_FINITE_PRJ2_DOMRES	$\text{finite}(E \triangleleft \text{prj}_2) \hat{=} \text{finite}(E)$	
* SIMP_FINITE_NATURAL	$\text{finite}(\mathbb{N}) \hat{=} \perp$	
* SIMP_FINITE_NATURAL1	$\text{finite}(\mathbb{N}_1) \hat{=} \perp$	
* SIMP_FINITE_INTEGER	$\text{finite}(\mathbb{Z}) \hat{=} \perp$	
* SIMP_FINITE_BOOL	$\text{finite}(\text{BOOL}) \hat{=} \top$	
* SIMP_FINITE_LAMBDA	$\text{finite}(\{x \cdot P \mid E \mapsto F\}) \hat{=} \text{finite}(\{x \cdot P \mid E\})$	where E maplet combinatic bound iden and expres that are not by the comprehen (i.e., E is syntactical injective) a identifiers by the comprehen that occur i also occur
* SIMP_TYPE_IN	$t \in Ty \hat{=} \top$	where Ty i expression
* SIMP_SPECIAL_EQV_BTRUE	$P \Leftrightarrow \top \hat{=} P$	
* SIMP_SPECIAL_EQV_BFALSE	$P \Leftrightarrow \perp \hat{=} \neg P$	
* DEF_SUBSET	$A \subset B \hat{=} A \subseteq B \wedge \neg A = B$	
* SIMP_SPECIAL_SUBSET_R	$S \subset \emptyset \hat{=} \perp$	
* SIMP_SPECIAL_SUBSET_L	$\emptyset \subset S \hat{=} \neg S = \emptyset$	
* SIMP_TYPE_SUBSET_L	$S \subset Ty \hat{=} S \neq Ty$	where Ty i expression
* SIMP_MULTI_SUBSET	$S \subset S \hat{=} \perp$	
* SIMP_EQUAL_CONSTR	$\text{constr}(a_1, \dots, a_n) = \text{constr}(b_1, \dots, b_n) \hat{=} a_1 = b_1 \wedge \dots \wedge a_n = b_n$	where CON datatype constructor
* SIMP_EQUAL_CONSTR_DIFF	$\text{constr}_1(\dots) = \text{constr}_2(\dots) \hat{=} \perp$	where CON and CONST different d constructor

*	SIMP_DESTR_CONSTR	$\text{destr}(\text{constr}(a_1, \dots, a_n)) \hat{=} a_i$	where des datatype d_i for the i -th argument c datatype constructor constr
*	DISTR1_AND_OR	$P \wedge (Q \vee R) \hat{=} (P \wedge Q) \vee (P \wedge R)$	
*	DISTR1_OR_AND	$P \vee (Q \wedge R) \hat{=} (P \vee Q) \wedge (P \vee R)$	
*	DEF_OR	$P \vee Q \vee \dots \vee R \hat{=} \neg P \Rightarrow (Q \vee \dots \vee R)$	
*	DERIV_IMP	$P \Rightarrow Q \hat{=} \neg Q \Rightarrow \neg P$	
*	DERIV_IMP_IMP	$P \Rightarrow (Q \Rightarrow R) \hat{=} P \wedge Q \Rightarrow R$	
*	DISTR1_IMP_AND	$P \Rightarrow (Q \wedge R) \hat{=} (P \Rightarrow Q) \wedge (P \Rightarrow R)$	
*	DISTR1_IMP_OR	$(P \vee Q) \Rightarrow R \hat{=} (P \Rightarrow R) \wedge (Q \Rightarrow R)$	
*	DEF_EQV	$P \Leftrightarrow Q \hat{=} (P \Rightarrow Q) \wedge (Q \Rightarrow P)$	
*	DISTR1_NOT_AND	$\neg(P \wedge Q) \hat{=} \neg P \vee \neg Q$	
*	DISTR1_NOT_OR	$\neg(P \vee Q) \hat{=} \neg P \wedge \neg Q$	
*	DERIV_NOT_IMP	$\neg(P \Rightarrow Q) \hat{=} P \wedge \neg Q$	
*	DERIV_NOT_FORALL	$\neg \forall x. P \hat{=} \exists x. \neg P$	
*	DERIV_NOT_EXISTS	$\neg \exists x. P \hat{=} \forall x. \neg P$	
*	DEF_IN_MAPSTO	$E \mapsto F \in S \times T \hat{=} E \in S \wedge F \in T$	
*	DEF_IN_POW	$E \in \mathbb{P}(S) \hat{=} E \subseteq S$	
*	DEF_IN_POW1	$E \in \mathbb{P}_1(S) \hat{=} E \in \mathbb{P}(S) \wedge S \neq \emptyset$	
*	DEF_SUBSETEQ	$S \subseteq T \hat{=} \forall x. x \in S \Rightarrow x \in T$	where x is in S or T
*	DEF_IN_BUNION	$E \in S \cup T \hat{=} E \in S \vee E \in T$	
*	DEF_IN_BINTER	$E \in S \cap T \hat{=} E \in S \wedge E \in T$	
*	DEF_IN_SETMINUS	$E \in S \setminus T \hat{=} E \in S \wedge \neg(E \in T)$	
*	DEF_IN_SETENUM	$E \in \{A, \dots, B\} \hat{=} E = A \vee \dots \vee E = B$	
*	DEF_IN_KUNION	$E \in \text{union}(S) \hat{=} \exists s. s \in S \wedge E \in s$	where s is
*	DEF_IN_QUNION	$E \in (\bigcup x. P(x) \mid T(x)) \hat{=} \exists s. P(s) \wedge E \in T(s)$	where s is
*	DEF_IN_KINTER	$E \in \text{inter}(S) \hat{=} \forall s. s \in S \Rightarrow E \in s$	where s is
*	DEF_IN_QINTER	$E \in (\bigcap x. P(x) \mid T(x)) \hat{=} \forall s. P(s) \Rightarrow E \in T(s)$	where s is
*	DEF_IN_UPTO	$E \in a .. b \hat{=} a \leq E \wedge E \leq b$	
*	DISTR1_BUNION_BINTER	$S \cup (T \cap U) \hat{=} (S \cup T) \cap (S \cup U)$	
*	DISTR1_BINTER_BUNION	$S \cap (T \cup U) \hat{=} (S \cap T) \cup (S \cap U)$	
	DISTR1_BINTER_SETMINUS	$S \cap (T \setminus U) \hat{=} (S \cap T) \setminus (S \cap U)$	
	DISTR1_SETMINUS_BUNION	$S \setminus (T \cup U) \hat{=} S \setminus T \setminus U$	
*	DERIV_TYPE_SETMINUS_BINTER	$Ty \setminus (S \cap T) \hat{=} (Ty \setminus S) \cup (Ty \setminus T)$	where Ty is expression
*	DERIV_TYPE_SETMINUS_BUNION	$Ty \setminus (S \cup T) \hat{=} (Ty \setminus S) \cap (Ty \setminus T)$	where Ty is expression
*	DERIV_TYPE_SETMINUS_SETMINUS	$Ty \setminus (S \setminus T) \hat{=} (Ty \setminus S) \cup T$	where Ty is expression
	DISTR1_CPROD_BINTER	$S \times (T \cap U) \hat{=} (S \times T) \cap (S \times U)$	
	DISTR1_CPROD_BUNION	$S \times (T \cup U) \hat{=} (S \times T) \cup (S \times U)$	
	DISTR1_CPROD_SETMINUS	$S \times (T \setminus U) \hat{=} (S \times T) \setminus (S \times U)$	
			where $\mathbb{P}(\cdot)$

*	DERIV_SUBSETEQ	$S \subseteq T \hat{=} (Ty \setminus T) \subseteq (Ty \setminus S)$	the type of T
*	DERIV_EQUAL	$S = T \hat{=} S \subseteq T \wedge T \subseteq S$	where $\mathbb{P}(\cdot)$ the type of T
*	DERIV_SUBSETEQ_SETMINUS_L	$A \setminus B \subseteq S \hat{=} A \subseteq B \cup S$	
*	DERIV_SUBSETEQ_SETMINUS_R	$S \subseteq A \setminus B \hat{=} S \subseteq A \wedge S \cap B = \emptyset$	
*	DEF_PARTITION	$\text{partition}(s, s_1, s_2, \dots, s_n) \hat{=} \begin{array}{l} s = s_1 \cup s_2 \cup \dots \cup s_n \\ \wedge s_1 \cap s_2 = \emptyset \\ \vdots \\ \wedge s_1 \cap s_n = \emptyset \\ \vdots \\ \wedge s_{n-1} \cap s_n = \emptyset \end{array}$	
	SIMP_EMPTY_PARTITION	$\text{partition}(S) \hat{=} S = \emptyset$	
	SIMP_SINGLE_PARTITION	$\text{partition}(S, T) \hat{=} S = T$	

Relation Rewrite Rules

Rules that are marked with a * in the first column are implemented in the latest version of Rodin. Rules without a * are planned to be implemented in future versions. Other conventions used in these tables are described in The_Proving_Perspective_(Rodin_User_Manual)#Rewrite_Rules.

	Name	Rule	Side Condi
*	SIMP_DOM_SETENUM	$\text{dom}(\{x \mapsto a, \dots, y \mapsto b\}) \hat{=} \{x, \dots, y\}$	
*	SIMP_DOM_CONVERSE	$\text{dom}(r^{-1}) \hat{=} \text{ran}(r)$	
*	SIMP_RAN_SETENUM	$\text{ran}(\{x \mapsto a, \dots, y \mapsto b\}) \hat{=} \{a, \dots, b\}$	
*	SIMP_RAN_CONVERSE	$\text{ran}(r^{-1}) \hat{=} \text{dom}(r)$	
*	SIMP_SPECIAL_OVERL	$r \triangleleft \dots \triangleleft \emptyset \triangleleft \dots \triangleleft s \hat{=} r \triangleleft \dots \triangleleft s$	
*	SIMP_MULTI_OVERL	$r_1 \triangleleft \dots \triangleleft r_n \hat{=} r_1 \triangleleft \dots \triangleleft r_{i-1} \triangleleft r_{i+1} \triangleleft \dots \triangleleft r_n$	there is a j such that $1 \leq i < j \leq n$ and i, j syntactically equal.
*	SIMP_TYPE_OVERL_CPROD	$r \triangleleft \dots \triangleleft Ty \triangleleft \dots \triangleleft s \hat{=} Ty \triangleleft \dots \triangleleft s$	where Ty is a type expr
*	SIMP_SPECIAL_DOMRES_L	$\emptyset \triangleleft r \hat{=} \emptyset$	
*	SIMP_SPECIAL_DOMRES_R	$S \triangleleft \emptyset \hat{=} \emptyset$	
*	SIMP_TYPE_DOMRES	$Ty \triangleleft r \hat{=} r$	where Ty is a type expr
*	SIMP_MULTI_DOMRES_DOM	$\text{dom}(r) \triangleleft r \hat{=} r$	
*	SIMP_MULTI_DOMRES_RAN	$\text{ran}(r) \triangleleft r^{-1} \hat{=} r^{-1}$	
*	SIMP_DOMRES_DOMRES_ID	$S \triangleleft (T \triangleleft \text{id}) \hat{=} (S \cap T) \triangleleft \text{id}$	
*	SIMP_DOMRES_DOMSUB_ID	$S \triangleleft (T \triangleleft \text{id}) \hat{=} (S \setminus T) \triangleleft \text{id}$	
*	SIMP_SPECIAL_RANRES_R	$r \triangleright \emptyset \hat{=} \emptyset$	
*	SIMP_SPECIAL_RANRES_L	$\emptyset \triangleright S \hat{=} \emptyset$	
*	SIMP_TYPE_RANRES	$r \triangleright Ty \hat{=} r$	where Ty is a type expr
*	SIMP_MULTI_RANRES_RAN	$r \triangleright \text{ran}(r) \hat{=} r$	
*	SIMP_MULTI_RANRES_DOM	$r^{-1} \triangleright \text{dom}(r) \hat{=} r^{-1}$	
*	SIMP_RANRES_ID	$\text{id} \triangleright S \hat{=} S \triangleleft \text{id}$	
*	SIMP_RANSUB_ID	$\text{id} \triangleright S \hat{=} S \triangleleft \text{id}$	
*	SIMP_RANRES_DOMRES_ID	$(S \triangleleft \text{id}) \triangleright T \hat{=} (S \cap T) \triangleleft \text{id}$	
*	SIMP_RANRES_DOMSUB_ID	$(S \triangleleft \text{id}) \triangleright T \hat{=} (T \setminus S) \triangleleft \text{id}$	
*	SIMP_SPECIAL_DOMSUB_L	$\emptyset \triangleleft r \hat{=} r$	

* SIMP_SPECIAL_DOMSUB_R	$S \triangleleft \emptyset \hat{=} \emptyset$	
* SIMP_TYPE_DOMSUB	$Ty \triangleleft r \hat{=} \emptyset$	where Ty is a type expr
* SIMP_MULTI_DOMSUB_DOM	$\text{dom}(r) \triangleleft r \hat{=} \emptyset$	
* SIMP_MULTI_DOMSUB_RAN	$\text{ran}(r) \triangleleft r^{-1} \hat{=} \emptyset$	
* SIMP_DOMSUB_DOMRES_ID	$S \triangleleft (T \triangleleft \text{id}) \hat{=} (T \setminus S) \triangleleft \text{id}$	
* SIMP_DOMSUB_DOMSUB_ID	$S \triangleleft (T \triangleleft \text{id}) \hat{=} (S \cup T) \triangleleft \text{id}$	
* SIMP_SPECIAL_RANSUB_R	$r \triangleright \emptyset \hat{=} r$	
* SIMP_SPECIAL_RANSUB_L	$\emptyset \triangleright S \hat{=} \emptyset$	
* SIMP_TYPE_RANSUB	$r \triangleright Ty \hat{=} \emptyset$	where Ty is a type expr
* SIMP_MULTI_RANSUB_DOM	$r^{-1} \triangleright \text{dom}(r) \hat{=} \emptyset$	
* SIMP_MULTI_RANSUB_RAN	$r \triangleright \text{ran}(r) \hat{=} \emptyset$	
* SIMP_RANSUB_DOMRES_ID	$(S \triangleleft \text{id}) \triangleright T \hat{=} (S \setminus T) \triangleleft \text{id}$	
* SIMP_RANSUB_DOMSUB_ID	$(S \triangleleft \text{id}) \triangleright T \hat{=} (S \cup T) \triangleleft \text{id}$	
* SIMP_SPECIAL_FCOMP	$r; \dots; \emptyset; \dots; s \hat{=} \emptyset$	
* SIMP_TYPE_FCOMP_ID	$r; \dots; \text{id}; \dots; s \hat{=} r; \dots; s$	
* SIMP_TYPE_FCOMP_R	$r; Ty \hat{=} \text{dom}(r) \times Tb$	where Ty is a type expr to $Ta \times Tb$
* SIMP_TYPE_FCOMP_L	$Ty; r \hat{=} Ta \times \text{ran}(r)$	where Ty is a type expr to $Ta \times Tb$
* SIMP_FCOMP_ID	$r; \dots; S \triangleleft \text{id}; T \triangleleft \text{id}; \dots; s \hat{=} r; \dots; (S \cap T) \triangleleft \text{id}; \dots; s$	
* SIMP_SPECIAL_BCOMP	$r \circ \dots \circ \emptyset \circ \dots \circ s \hat{=} \emptyset$	
* SIMP_TYPE_BCOMP_ID	$r \circ \dots \circ \text{id} \circ \dots \circ s \hat{=} r \circ \dots \circ s$	
* SIMP_TYPE_BCOMP_L	$Ty \circ r \hat{=} \text{dom}(r) \times Tb$	where Ty is a type expr to $Ta \times Tb$
* SIMP_TYPE_BCOMP_R	$r \circ Ty \hat{=} Ta \times \text{ran}(r)$	where Ty is a type expr to $Ta \times Tb$
* SIMP_BCOMP_ID	$r \circ \dots \circ S \triangleleft \text{id} \circ T \triangleleft \text{id} \circ \dots \circ s \hat{=} r \circ \dots \circ (S \cap T) \triangleleft \text{id} \circ \dots \circ s$	
* SIMP_SPECIAL_DPROD_R	$r \otimes \emptyset \hat{=} \emptyset$	
* SIMP_SPECIAL_DPROD_L	$\emptyset \otimes r \hat{=} \emptyset$	
* SIMP_DPROD_CPROD	$(S \times T) \otimes (U \times V) \hat{=} S \cap U \times (T \times V)$	
* SIMP_SPECIAL_PPROD_R	$r \parallel \emptyset \hat{=} \emptyset$	
* SIMP_SPECIAL_PPROD_L	$\emptyset \parallel r \hat{=} \emptyset$	
* SIMP_PPROD_CPROD	$(S \times T) \parallel (U \times V) \hat{=} (S \times U) \times (T \times V)$	
* SIMP_SPECIAL_RELIMAGE_R	$r[\emptyset] \hat{=} \emptyset$	
* SIMP_SPECIAL_RELIMAGE_L	$\emptyset[S] \hat{=} \emptyset$	
* SIMP_TYPE_RELIMAGE	$r[Ty] \hat{=} \text{ran}(r)$	where Ty is a type expr
* SIMP_MULTI_RELIMAGE_DOM	$r[\text{dom}(r)] \hat{=} \text{ran}(r)$	
* SIMP_RELIMAGE_ID	$\text{id}[T] \hat{=} T$	
* SIMP_RELIMAGE_DOMRES_ID	$(S \triangleleft \text{id})[T] \hat{=} S \cap T$	
* SIMP_RELIMAGE_DOMSUB_ID	$(S \triangleleft \text{id})[T] \hat{=} T \setminus S$	
* SIMP_MULTI_RELIMAGE_CPROD_SING	$(\{E\} \times S)[\{E\}] \hat{=} S$	where E is a single ex
* SIMP_MULTI_RELIMAGE_SING_MAPSTO	$\{E \mapsto F\}[\{E\}] \hat{=} \{F\}$	where E is a single ex
* SIMP_MULTI_RELIMAGE_CONVERSE_RANSUB	$(r \triangleright S)^{-1}[S] \hat{=} \emptyset$	
* SIMP_MULTI_RELIMAGE_CONVERSE_RANRES	$(r \triangleright S)^{-1}[S] \hat{=} r^{-1}[S]$	
* SIMP_RELIMAGE_CONVERSE_DOMSUB	$(S \triangleleft r)^{-1}[T] \hat{=} r^{-1}[T] \setminus S$	

	DERIV_RELIMAGE_RANSUB	$(r \triangleright S)[T] \hat{=} r[T] \setminus S$	
	DERIV_RELIMAGE_RANRES	$(r \triangleright S)[T] \hat{=} r[T] \cap S$	
*	SIMP_MULTI_RELIMAGE_DOMSUB	$(S \triangleleft r)[S] \hat{=} \emptyset$	
	DERIV_RELIMAGE_DOMSUB	$(S \triangleleft r)[T] \hat{=} r[T \setminus S]$	
	DERIV_RELIMAGE_DOMRES	$(S \triangleleft r)[T] \hat{=} r[S \cap T]$	
*	SIMP_SPECIAL_CONVERSE	$\emptyset^{-1} \hat{=} \emptyset$	
*	SIMP_CONVERSE_ID	$\text{id}^{-1} \hat{=} \text{id}$	
*	SIMP_CONVERSE_CPROD	$(S \times T)^{-1} \hat{=} T \times S$	
*	SIMP_CONVERSE_SETENUM	$\{x \mapsto a, \dots, y \mapsto b\}^{-1} \hat{=} \{a \mapsto x, \dots, b \mapsto y\}$	
*	SIMP_CONVERSE_COMPSET	$\{X \cdot P \mid x \mapsto y\}^{-1} \hat{=} \{X \cdot P \mid y \mapsto x\}$	
*	SIMP_DOM_ID	$\text{dom}(\text{id}) \hat{=} S$	where id has type $\mathbb{P}(S)$
*	SIMP_RAN_ID	$\text{ran}(\text{id}) \hat{=} S$	where id has type $\mathbb{P}(S)$
*	SIMP_FCOMP_ID_L	$(S \triangleleft \text{id}); r \hat{=} S \triangleleft r$	
*	SIMP_FCOMP_ID_R	$r; (S \triangleleft \text{id}) \hat{=} r \triangleright S$	
*	SIMP_SPECIAL_REL_R	$S \leftrightarrow \emptyset \hat{=} \{\emptyset\}$	idem for operators \leftrightarrow
*	SIMP_SPECIAL_REL_L	$\emptyset \leftrightarrow S \hat{=} \{\emptyset\}$	idem for operators $\leftrightarrow \mapsto \rightarrow \mapsto \mapsto$
*	SIMP_FUNIMAGE_PRJ1	$\text{prj}_1(E \mapsto F) \hat{=} E$	
*	SIMP_FUNIMAGE_PRJ2	$\text{prj}_2(E \mapsto F) \hat{=} F$	
*	SIMP_DOM_PRJ1	$\text{dom}(\text{prj}_1) \hat{=} S \times T$	where prj_1 has type $\mathbb{P}(S \times T \times S)$
*	SIMP_DOM_PRJ2	$\text{dom}(\text{prj}_2) \hat{=} S \times T$	where prj_2 has type $\mathbb{P}(S \times T \times T)$
*	SIMP_RAN_PRJ1	$\text{ran}(\text{prj}_1) \hat{=} S$	where prj_1 has type $\mathbb{P}(S \times T \times S)$
*	SIMP_RAN_PRJ2	$\text{ran}(\text{prj}_2) \hat{=} T$	where prj_2 has type $\mathbb{P}(S \times T \times T)$
*	SIMP_FUNIMAGE_LAMBDA	$(\lambda x \cdot P(x) \mid E(x))(y) \hat{=} E(y)$	
*	SIMP_DOM_LAMBDA	$\text{dom}(\{x \cdot P \mid E \mapsto F\}) \hat{=} \{x \cdot P \mid E\}$	
*	SIMP_RAN_LAMBDA	$\text{ran}(\{x \cdot P \mid E \mapsto F\}) \hat{=} \{x \cdot P \mid F\}$	
*	SIMP_IN_FUNIMAGE	$E \mapsto F(E) \in F \hat{=} \top$	
*	SIMP_IN_FUNIMAGE_CONVERSE_L	$F^{-1}(E) \mapsto E \in F \hat{=} \top$	
*	SIMP_IN_FUNIMAGE_CONVERSE_R	$F(E) \mapsto E \in F^{-1} \hat{=} \top$	
*	SIMP_MULTI_FUNIMAGE_SETENUM_LL	$\{A \mapsto E, \dots, B \mapsto E\}(x) \hat{=} E$	
*	SIMP_MULTI_FUNIMAGE_SETENUM_LR	$\{E, \dots, x \mapsto y, \dots, F\}(x) \hat{=} y$	
*	SIMP_MULTI_FUNIMAGE_OVERL_SETENUM	$(r \triangleleft \dots \triangleleft \{E, \dots, x \mapsto y, \dots, F\})(x) \hat{=} y$	
*	SIMP_MULTI_FUNIMAGE_BUNION_SETENUM	$(r \cup \dots \cup \{E, \dots, x \mapsto y, \dots, F\})(x) \hat{=} y$	
*	SIMP_FUNIMAGE_CPROD	$(S \times \{F\})(x) \hat{=} F$	
*	SIMP_FUNIMAGE_ID	$\text{id}(x) \hat{=} x$	
*	SIMP_FUNIMAGE_FUNIMAGE_CONVERSE	$f(f^{-1}(E)) \hat{=} E$	
*	SIMP_FUNIMAGE_CONVERSE_FUNIMAGE	$f^{-1}(f(E)) \hat{=} E$	
*	SIMP_FUNIMAGE_FUNIMAGE_CONVERSE_SETENUM	$\{x \mapsto a, \dots, y \mapsto b\}(\{a \mapsto x, \dots, b \mapsto y\}(E)) \hat{=} E$	
*	SIMP_FUNIMAGE_DOMRES	$(E \triangleleft F)(G) \hat{=} F(G)$	with hypothesis $F \in A$ where op is one of $\mapsto, \mapsto, \mapsto$.

*	SIMP_FUNIMAGE_DOMSUB	$(E \triangleleft F)(G) \hat{=} F(G)$	with hypothesis $F \in A$ where op is one of \mapsto , $\mapsto\mapsto$, $\mapsto\mapsto\mapsto$.
*	SIMP_FUNIMAGE_RANRES	$(F \triangleright E)(G) \hat{=} F(G)$	with hypothesis $F \in A$ where op is one of \mapsto , $\mapsto\mapsto$, $\mapsto\mapsto\mapsto$.
*	SIMP_FUNIMAGE_RANSUB	$(F \triangleright E)(G) \hat{=} F(G)$	with hypothesis $F \in A$ where op is one of \mapsto , $\mapsto\mapsto$, $\mapsto\mapsto\mapsto$.
*	SIMP_FUNIMAGE_SETMINUS	$(F \setminus E)(G) \hat{=} F(G)$	with hypothesis $F \in A$ where op is one of \mapsto , $\mapsto\mapsto$, $\mapsto\mapsto\mapsto$.
	DEF_BCOMP	$r \circ \dots \circ s \hat{=} s; \dots; r$	
	DERIV_ID_SING	$\{E\} \triangleleft id \hat{=} \{E \mapsto E\}$	where E is a single ex
*	SIMP_SPECIAL_DOM	$dom(\emptyset) \hat{=} \emptyset$	
*	SIMP_SPECIAL_RAN	$ran(\emptyset) \hat{=} \emptyset$	
*	SIMP_CONVERSE_CONVERSE	$r^{-1-1} \hat{=} r$	
*	DERIV_RELIMAGE_FCOMP	$(p; q)[s] \hat{=} q[p[s]]$	
*	DERIV_FCOMP_DOMRES	$(s \triangleleft p); q \hat{=} s \triangleleft (p; q)$	
*	DERIV_FCOMP_DOMSUB	$(s \triangleleft p); q \hat{=} s \triangleleft (p; q)$	
*	DERIV_FCOMP_RANRES	$p; (q \triangleright s) \hat{=} (p; q) \triangleright s$	
*	DERIV_FCOMP_RANSUB	$p; (q \triangleright s) \hat{=} (p; q) \triangleright s$	
	DERIV_FCOMP_SING	$\{E \mapsto F\}; \{F \mapsto G\} \hat{=} \{E \mapsto G\}$	
*	SIMP_SPECIAL_EQUAL_RELDOMRAN	$\emptyset \leftrightarrow \emptyset \hat{=} \{\emptyset\}$	idem for operators $\mapsto\mapsto$
*	SIMP_TYPE_DOM	$dom(Ty) \hat{=} Ta$	where Ty is a type expr to $Ta \times Tb$
*	SIMP_TYPE_RAN	$ran(Ty) \hat{=} Tb$	where Ty is a type expr to $Ta \times Tb$
*	SIMP_MULTI_DOM_CPROD	$dom(E \times E) \hat{=} E$	
*	SIMP_MULTI_RAN_CPROD	$ran(E \times E) \hat{=} E$	
*	SIMP_MULTI_DOM_DOMRES	$dom(A \triangleleft f) \hat{=} dom(f) \cap A$	
*	SIMP_MULTI_DOM_DOMSUB	$dom(A \triangleleft f) \hat{=} dom(f) \setminus A$	
*	SIMP_MULTI_RAN_RANRES	$ran(f \triangleright A) \hat{=} ran(f) \cap A$	
*	SIMP_MULTI_RAN_RANSUB	$ran(f \triangleright A) \hat{=} ran(f) \setminus A$	
*	DEF_IN_DOM	$E \in dom(r) \hat{=} \exists y. E \mapsto y \in r$	
*	DEF_IN_RAN	$F \in ran(r) \hat{=} \exists x. x \mapsto F \in r$	
*	DEF_IN_CONVERSE	$E \mapsto F \in r^{-1} \hat{=} F \mapsto E \in r$	
*	DEF_IN_DOMRES	$E \mapsto F \in S \triangleleft r \hat{=} E \in S \wedge E \mapsto F \in r$	
*	DEF_IN_RANRES	$E \mapsto F \in r \triangleright T \hat{=} E \mapsto F \in r \wedge F \in T$	
*	DEF_IN_DOMSUB	$E \mapsto F \in S \triangleleft r \hat{=} E \notin S \wedge E \mapsto F \in r$	
*	DEF_IN_RANSUB	$E \mapsto F \in r \triangleright T \hat{=} E \mapsto F \in r \wedge F \notin T$	
*	DEF_IN_RELIMAGE	$F \in r[w] \hat{=} \exists x. x \in w \wedge x \mapsto F \in r$	
*	DEF_IN_FCOMP	$E \mapsto F \in (p; q) \hat{=} \exists x. E \mapsto x \in p \wedge x \mapsto F \in q$	
*	DEF_OVERL	$p \triangleleft q \hat{=} (dom(q) \triangleleft p) \cup q$	
*	DEF_IN_ID	$E \mapsto F \in id \hat{=} E = F$	
*	DEF_IN_DPROD	$E \mapsto (F \mapsto G) \in p \otimes q \hat{=} E \mapsto F \in p \wedge E \mapsto G \in q$	
*	DEF_IN_PPROD	$(E \mapsto G) \mapsto (F \mapsto H) \in p \parallel q \hat{=} E \mapsto F \in p \wedge G \mapsto H \in q$	
*	DEF_IN_REL	$r \in S \leftrightarrow T \hat{=} r \subseteq S \times T$	

*	DEF_IN_RELDOM	$r \in S \leftrightarrow T \hat{=} r \in S \leftrightarrow T \wedge \text{dom}(r) = S$	
*	DEF_IN_RELRAN	$r \in S \leftrightarrow T \hat{=} r \in S \leftrightarrow T \wedge \text{ran}(r) = T$	
*	DEF_IN_RELDOMRAN	$r \in S \leftrightarrow T \hat{=} r \in S \leftrightarrow T \wedge \text{dom}(r) = S \wedge \text{ran}(r) = T$	
*	DEF_IN_FCT	$f \in S \rightarrow T \hat{=} f \in S \leftrightarrow T$ $\wedge (\forall x, y, z. x \mapsto y \in f \wedge x \mapsto z \in f \Rightarrow y = z)$	
*	DEF_IN_TFCT	$f \in S \rightarrow T \hat{=} f \in S \leftrightarrow T \wedge \text{dom}(f) = S$	
*	DEF_IN_INJ	$f \in S \rightarrow T \hat{=} f \in S \leftrightarrow T \wedge f^{-1} \in T \rightarrow S$	
*	DEF_IN_TINJ	$f \in S \rightarrow T \hat{=} f \in S \rightarrow T \wedge \text{dom}(f) = S$	
*	DEF_IN_SURJ	$f \in S \rightarrow T \hat{=} f \in S \leftrightarrow T \wedge \text{ran}(f) = T$	
*	DEF_IN_TSURJ	$f \in S \rightarrow T \hat{=} f \in S \rightarrow T \wedge \text{dom}(f) = S$	
*	DEF_IN_BIJ	$f \in S \rightarrow T \hat{=} f \in S \rightarrow T \wedge \text{ran}(f) = T$	
	DISTRIBCOMP_BUNION	$r \circ (s \cup t) \hat{=} (r \circ s) \cup (r \circ t)$	
*	DISTRIBCOMP_BUNION_R	$p; (q \cup r) \hat{=} (p; q) \cup (p; r)$	
*	DISTRIBCOMP_BUNION_L	$(q \cup r); p \hat{=} (q; p) \cup (r; p)$	
	DISTRIBPROD_BUNION	$r \otimes (s \cup t) \hat{=} (r \otimes s) \cup (r \otimes t)$	
	DISTRIBPROD_BINTER	$r \otimes (s \cap t) \hat{=} (r \otimes s) \cap (r \otimes t)$	
	DISTRIBPROD_SETMINUS	$r \otimes (s \setminus t) \hat{=} (r \otimes s) \setminus (r \otimes t)$	
	DISTRIBPROD_OVERL	$r \otimes (s \Leftarrow t) \hat{=} (r \otimes s) \Leftarrow (r \otimes t)$	
	DISTRIBPROD_BUNION	$r \parallel (s \cup t) \hat{=} (r \parallel s) \cup (r \parallel t)$	
	DISTRIBPROD_BINTER	$r \parallel (s \cap t) \hat{=} (r \parallel s) \cap (r \parallel t)$	
	DISTRIBPROD_SETMINUS	$r \parallel (s \setminus t) \hat{=} (r \parallel s) \setminus (r \parallel t)$	
	DISTRIBPROD_OVERL	$r \parallel (s \Leftarrow t) \hat{=} (r \parallel s) \Leftarrow (r \parallel t)$	
	DISTRIBOVERL_BUNION_L	$(p \cup q) \Leftarrow r \hat{=} (p \Leftarrow r) \cup (q \Leftarrow r)$	
	DISTRIBOVERL_BINTER_L	$(p \cap q) \Leftarrow r \hat{=} (p \Leftarrow r) \cap (q \Leftarrow r)$	
*	DISTRIBDOMRES_BUNION_R	$s \triangleleft (p \cup q) \hat{=} (s \triangleleft p) \cup (s \triangleleft q)$	
*	DISTRIBDOMRES_BUNION_L	$(s \cup t) \triangleleft r \hat{=} (s \triangleleft r) \cup (t \triangleleft r)$	
*	DISTRIBDOMRES_BINTER_R	$s \triangleleft (p \cap q) \hat{=} (s \triangleleft p) \cap (s \triangleleft q)$	
*	DISTRIBDOMRES_BINTER_L	$(s \cap t) \triangleleft r \hat{=} (s \triangleleft r) \cap (t \triangleleft r)$	
	DISTRIBDOMRES_SETMINUS_R	$s \triangleleft (p \setminus q) \hat{=} (s \triangleleft p) \setminus (s \triangleleft q)$	
	DISTRIBDOMRES_SETMINUS_L	$(s \setminus t) \triangleleft r \hat{=} (s \triangleleft r) \setminus (t \triangleleft r)$	
	DISTRIBDOMRES_DPROD	$s \triangleleft (p \otimes q) \hat{=} (s \triangleleft p) \otimes (s \triangleleft q)$	
	DISTRIBDOMRES_OVERL	$s \triangleleft (r \Leftarrow q) \hat{=} (s \triangleleft r) \Leftarrow (s \triangleleft q)$	
*	DISTRIBDOMSUB_BUNION_R	$s \triangleleft (p \cup q) \hat{=} (s \triangleleft p) \cup (s \triangleleft q)$	
*	DISTRIBDOMSUB_BUNION_L	$(s \cup t) \triangleleft r \hat{=} (s \triangleleft r) \cap (t \triangleleft r)$	
*	DISTRIBDOMSUB_BINTER_R	$s \triangleleft (p \cap q) \hat{=} (s \triangleleft p) \cap (s \triangleleft q)$	
*	DISTRIBDOMSUB_BINTER_L	$(s \cap t) \triangleleft r \hat{=} (s \triangleleft r) \cup (t \triangleleft r)$	
	DISTRIBDOMSUB_DPROD	$A \triangleleft (r \otimes s) \hat{=} (A \triangleleft r) \otimes (A \triangleleft s)$	
	DISTRIBDOMSUB_OVERL	$A \triangleleft (r \Leftarrow s) \hat{=} (A \triangleleft r) \Leftarrow (A \triangleleft s)$	
*	DISTRIBRANRES_BUNION_R	$r \triangleright (s \cup t) \hat{=} (r \triangleright s) \cup (r \triangleright t)$	
*	DISTRIBRANRES_BUNION_L	$(p \cup q) \triangleright s \hat{=} (p \triangleright s) \cup (q \triangleright s)$	
*	DISTRIBRANRES_BINTER_R	$r \triangleright (s \cap t) \hat{=} (r \triangleright s) \cap (r \triangleright t)$	
*	DISTRIBRANRES_BINTER_L	$(p \cap q) \triangleright s \hat{=} (p \triangleright s) \cap (q \triangleright s)$	

	DISTR1_RANRES_SETMINUS_R	$r \triangleright (s \setminus t) \hat{=} (r \triangleright s) \setminus (r \triangleright t)$	
	DISTR1_RANRES_SETMINUS_L	$(p \setminus q) \triangleright s \hat{=} (p \triangleright s) \setminus (q \triangleright s)$	
*	DISTR1_RANSUB_BUNION_R	$r \triangleright (s \cup t) \hat{=} (r \triangleright s) \cap (r \triangleright t)$	
*	DISTR1_RANSUB_BUNION_L	$(p \cup q) \triangleright s \hat{=} (p \triangleright s) \cup (q \triangleright s)$	
*	DISTR1_RANSUB_BINTER_R	$r \triangleright (s \cap t) \hat{=} (r \triangleright s) \cup (r \triangleright t)$	
*	DISTR1_RANSUB_BINTER_L	$(p \cap q) \triangleright s \hat{=} (p \triangleright s) \cap (q \triangleright s)$	
*	DISTR1_CONVERSE_BUNION	$(p \cup q)^{-1} \hat{=} p^{-1} \cup q^{-1}$	
	DISTR1_CONVERSE_BINTER	$(p \cap q)^{-1} \hat{=} p^{-1} \cap q^{-1}$	
	DISTR1_CONVERSE_SETMINUS	$(r \setminus s)^{-1} \hat{=} r^{-1} \setminus s^{-1}$	
	DISTR1_CONVERSE_BCOMP	$(r \circ s)^{-1} \hat{=} (s^{-1} \circ r^{-1})$	
	DISTR1_CONVERSE_FCOMP	$(p ; q)^{-1} \hat{=} (q^{-1} ; p^{-1})$	
	DISTR1_CONVERSE_PPROD	$(r \parallel s)^{-1} \hat{=} r^{-1} \parallel s^{-1}$	
	DISTR1_CONVERSE_DOMRES	$(s \triangleleft r)^{-1} \hat{=} r^{-1} \triangleright s$	
	DISTR1_CONVERSE_DOMSUB	$(s \triangleleft r)^{-1} \hat{=} r^{-1} \triangleright s$	
	DISTR1_CONVERSE_RANRES	$(r \triangleright s)^{-1} \hat{=} s \triangleleft r^{-1}$	
	DISTR1_CONVERSE_RANSUB	$(r \triangleright s)^{-1} \hat{=} s \triangleleft r^{-1}$	
*	DISTR1_DOM_BUNION	$\text{dom}(r \cup s) \hat{=} \text{dom}(r) \cup \text{dom}(s)$	
*	DISTR1_RAN_BUNION	$\text{ran}(r \cup s) \hat{=} \text{ran}(r) \cup \text{ran}(s)$	
*	DISTR1_RELIMAGE_BUNION_R	$r[S \cup T] \hat{=} r[S] \cup r[T]$	
*	DISTR1_RELIMAGE_BUNION_L	$(p \cup q)[S] \hat{=} p[S] \cup q[S]$	
*	DERIV_DOM_TOTALREL	$\text{dom}(r) \hat{=} E$	with hypothesis $r \in E$ where op is one of $\leftrightarrow, \leftrightarrow, \rightarrow, \mapsto, \twoheadrightarrow,$
	DERIV_RAN_SURJREL	$\text{ran}(r) \hat{=} F$	with hypothesis $r \in E$ where op is one of $\leftrightarrow, \leftrightarrow, \twoheadrightarrow, \twoheadrightarrow, \twoheadrightarrow$
*	DERIV_PRJ1_SURJ	$\text{prj}_1 \in Ty_1 \text{ op } Ty_2 \hat{=} \top$	where Ty_1 and Ty_2 are op is one of $\leftrightarrow, \leftrightarrow, \leftrightarrow, \leftrightarrow, \twoheadrightarrow,$
*	DERIV_PRJ2_SURJ	$\text{prj}_2 \in Ty_1 \text{ op } Ty_2 \hat{=} \top$	where Ty_1 and Ty_2 are op is one of $\leftrightarrow, \leftrightarrow, \leftrightarrow, \leftrightarrow, \twoheadrightarrow,$
*	DERIV_ID_BIJ	$\text{id} \in Ty \text{ op } Ty \hat{=} \top$	where Ty is a type and arrow
*	SIMP_MAPSTO_PRJ1_PRJ2	$\text{prj}_1(E) \mapsto \text{prj}_2(E) \hat{=} E$	
	DERIV_EXPAND_PRJS	$E \hat{=} \text{prj}_1(E) \mapsto \text{prj}_2(E)$	
*	SIMP_DOM_SUCC	$\text{dom}(\text{succ}) \hat{=} \mathbb{Z}$	
*	SIMP_RAN_SUCC	$\text{ran}(\text{succ}) \hat{=} \mathbb{Z}$	
*	DERIV_MULTI_IN_BUNION	$E \in A \cup \dots \cup \{\dots, E, \dots\} \cup \dots \cup B \hat{=} \top$	
*	DERIV_MULTI_IN_SETMINUS	$E \in S \setminus \{\dots, E, \dots\} \hat{=} \perp$	
*	DEF_PRED	$\text{pred} \hat{=} \text{succ}^{-1}$	

Empty Set Rewrite Rules

Rules that are marked with a * in the first column are implemented in the latest version of Rodin. Rules without a * are planned to be implemented in future versions. Other conventions used in these tables are described in [The_Proving_Perspective_\(Rodin_User_Manual\)#Rewrite_Rules](#).

All rewrite rules that match the pattern $\mathbf{P} = \emptyset$ are also applicable to predicates of the form $\mathbf{P} \subseteq \emptyset$ and $\emptyset = \mathbf{P}$, as these predicates are equivalent. All rewrite rules that match the pattern $\mathbf{P} = Ty$ are also applicable to predicates of the form $Ty \subseteq \mathbf{P}$ and $Ty = \mathbf{P}$, as these predicates are equivalent.

	Name	Rule	Side Condition
*	DEF_SPECIAL_NOT_EQUAL	$\neg S = \emptyset \hat{=} \exists x \cdot x \in S$	where x is not free in
*	SIMP_SETENUM_EQUAL_EMPTY	$\{A, \dots, B\} = \emptyset \hat{=} \perp$	
*	SIMP_SPECIAL_EQUAL_COMPSET	$\{x \cdot P(x) \mid E\} = \emptyset \hat{=} \forall x \cdot \neg P(x)$	
*	SIMP_BINTER_EQUAL_TYPE	$A \cap \dots \cap B = Ty \hat{=} A = Ty \wedge \dots \wedge B = Ty$	where Ty is a type exp
*	SIMP_BINTER_SING_EQUAL_EMPTY	$A \cap \dots \cap \{a\} \cap \dots \cap B = \emptyset \hat{=} \neg a \in A \cap \dots \cap B$	
*	SIMP_BINTER_SETMINUS_EQUAL_EMPTY	$A \cap \dots \cap (B \setminus C) \cap \dots \cap D = \emptyset \hat{=} (A \cap \dots \cap B \cap \dots \cap D) \setminus C = \emptyset$	
*	SIMP_BUNION_EQUAL_EMPTY	$A \cup \dots \cup B = \emptyset \hat{=} A = \emptyset \wedge \dots \wedge B = \emptyset$	
*	SIMP_SETMINUS_EQUAL_EMPTY	$A \setminus B = \emptyset \hat{=} A \subseteq B$	
*	SIMP_SETMINUS_EQUAL_TYPE	$A \setminus B = Ty \hat{=} A = Ty \wedge B = \emptyset$	where Ty is a type exp
*	SIMP_POW_EQUAL_EMPTY	$\mathbb{P}(S) = \emptyset \hat{=} \perp$	
*	SIMP_POW1_EQUAL_EMPTY	$\mathbb{P}_1(S) = \emptyset \hat{=} S = \emptyset$	
*	SIMP_KINTER_EQUAL_TYPE	$\text{inter}(S) = Ty \hat{=} S = \{Ty\}$	where Ty is a type exp
*	SIMP_KUNION_EQUAL_EMPTY	$\text{union}(S) = \emptyset \hat{=} S \subseteq \{\emptyset\}$	
*	SIMP_QINTER_EQUAL_TYPE	$(\bigcap x \cdot P(x) \mid E(x)) = Ty \hat{=} \forall x \cdot P(x) \Rightarrow E(x) = Ty$	where Ty is a type exp
*	SIMP_QUNION_EQUAL_EMPTY	$(\bigcup x \cdot P(x) \mid E(x)) = \emptyset \hat{=} \forall x \cdot P(x) \Rightarrow E(x) = \emptyset$	
*	SIMP_NATURAL_EQUAL_EMPTY	$\mathbb{N} = \emptyset \hat{=} \perp$	
*	SIMP_NATURAL1_EQUAL_EMPTY	$\mathbb{N}_1 = \emptyset \hat{=} \perp$	
*	SIMP_TYPE_EQUAL_EMPTY	$Ty = \emptyset \hat{=} \perp$	where Ty is a type exp
*	SIMP_CPROD_EQUAL_EMPTY	$S \times T = \emptyset \hat{=} S = \emptyset \vee T = \emptyset$	
*	SIMP_CPROD_EQUAL_TYPE	$S \times T = Ty \hat{=} S = Ta \wedge T = Tb$	where Ty is a type exp equal to $Ta \times Tb$
*	SIMP_UPTO_EQUAL_EMPTY	$i .. j = \emptyset \hat{=} i > j$	
*	SIMP_UPTO_EQUAL_INTEGER	$i .. j = \mathbb{Z} \hat{=} \perp$	
*	SIMP_UPTO_EQUAL_NATURAL	$i .. j = \mathbb{N} \hat{=} \perp$	
*	SIMP_UPTO_EQUAL_NATURAL1	$i .. j = \mathbb{N}_1 \hat{=} \perp$	
*	SIMP_SPECIAL_EQUAL_REL	$A \leftrightarrow B = \emptyset \hat{=} \perp$	idem for operators \leftrightarrow
	SIMP_TYPE_EQUAL_REL	$A \leftrightarrow B = Ty \hat{=} A = Ta \wedge B = Tb$	where Ty is a type exp equal to $Ta \times Tb$
*	SIMP_SPECIAL_EQUAL_RELDOM	$A \leftrightarrow B = \emptyset \hat{=} \neg A = \emptyset \wedge B = \emptyset$	idem for operator \rightarrow
	SIMP_TYPE_EQUAL_RELDOMRAN	$A \leftrightarrow B = Ty \hat{=} \perp$	where Ty is a type exp idem for operator $\leftrightarrow, \leftarrow, \rightarrow, \rightrightarrows, \rightleftarrows$
*	SIMP_SREL_EQUAL_EMPTY	$A \leftrightarrow B = \emptyset \hat{=} A = \emptyset \wedge \neg B = \emptyset$	
*	SIMP_STREL_EQUAL_EMPTY	$A \leftrightarrow B = \emptyset \hat{=} (A = \emptyset \Leftrightarrow \neg B = \emptyset)$	
*	SIMP_DOM_EQUAL_EMPTY	$\text{dom}(r) = \emptyset \hat{=} r = \emptyset$	
*	SIMP_RAN_EQUAL_EMPTY	$\text{ran}(r) = \emptyset \hat{=} r = \emptyset$	
*	SIMP_FCOMP_EQUAL_EMPTY	$p ; q = \emptyset \hat{=} \text{ran}(p) \cap \text{dom}(q) = \emptyset$	
*	SIMP_BCOMP_EQUAL_EMPTY	$p \circ q = \emptyset \hat{=} \text{ran}(q) \cap \text{dom}(p) = \emptyset$	
*	SIMP_DOMRES_EQUAL_EMPTY	$S \triangleleft r = \emptyset \hat{=} \text{dom}(r) \cap S = \emptyset$	
*	SIMP_DOMRES_EQUAL_TYPE	$S \triangleleft r = Ty \hat{=} S = Ta \wedge r = Ty$	where Ty is a type exp equal to $Ta \times Tb$
*	SIMP_DOMSUB_EQUAL_EMPTY	$S \triangleleft r = \emptyset \hat{=} \text{dom}(r) \subseteq S$	
*	SIMP_DOMSUB_EQUAL_TYPE	$S \triangleleft r = Ty \hat{=} S = \emptyset \wedge r = Ty$	where Ty is a type exp
*	SIMP_RANRES_EQUAL_EMPTY	$r \triangleright S = \emptyset \hat{=} \text{ran}(r) \cap S = \emptyset$	

*	SIMP_RANRES_EQUAL_TYPE	$r \triangleright S = Ty \hat{=} S = Tb \wedge r = Ty$	where Ty is a type exp equal to $Ta \times Tb$
*	SIMP_RANSUB_EQUAL_EMPTY	$r \triangleright S = \emptyset \hat{=} \text{ran}(r) \subseteq S$	
*	SIMP_RANSUB_EQUAL_TYPE	$r \triangleright S = Ty \hat{=} S = \emptyset \wedge r = Ty$	where Ty is a type exp
*	SIMP_CONVERSE_EQUAL_EMPTY	$r^{-1} = \emptyset \hat{=} r = \emptyset$	
*	SIMP_CONVERSE_EQUAL_TYPE	$r^{-1} = Ty \hat{=} r = Ty^{-1}$	where Ty is a type exp
*	SIMP_RELIMAGE_EQUAL_EMPTY	$r[S] = \emptyset \hat{=} S \triangleleft r = \emptyset$	
*	SIMP_OVERL_EQUAL_EMPTY	$r \triangleleft \dots \triangleleft s = \emptyset \hat{=} r = \emptyset \wedge \dots \wedge s = \emptyset$	
*	SIMP_DPROD_EQUAL_EMPTY	$p \otimes q = \emptyset \hat{=} \text{dom}(p) \cap \text{dom}(q) = \emptyset$	
*	SIMP_DPROD_EQUAL_TYPE	$p \otimes q = Ty \hat{=} p = Ta \times Tb \wedge q = Ta \times Tc$	where Ty is a type exp equal to $Ta \times (Tb \times Tc)$
*	SIMP_PPROD_EQUAL_EMPTY	$p \parallel q = \emptyset \hat{=} p = \emptyset \vee q = \emptyset$	
*	SIMP_PPROD_EQUAL_TYPE	$p \parallel q = Ty \hat{=} p = Ta \times Tc \wedge q = Tb \times Td$	where Ty is a type exp equal to $(Ta \times Tb) \times (Tc \times Td)$
*	SIMP_ID_EQUAL_EMPTY	$\text{id} = \emptyset \hat{=} \perp$	
*	SIMP_PRJ1_EQUAL_EMPTY	$\text{prj}_1 = \emptyset \hat{=} \perp$	
*	SIMP_PRJ2_EQUAL_EMPTY	$\text{prj}_2 = \emptyset \hat{=} \perp$	

Arithmetic Rewrite Rules

Rules that are marked with a * in the first column are implemented in the latest version of Rodin. Rules without a * are planned to be implemented in future versions. Other conventions used in these tables are described in [The_Proving_Perspective_\(Rodin_User_Manual\)#Rewrite_Rules](#).

	Name	Rule	Side Condition
*	SIMP_SPECIAL_MOD_0	$0 \bmod E \hat{=} 0$	
*	SIMP_SPECIAL_MOD_1	$E \bmod 1 \hat{=} 0$	
*	SIMP_MIN_SING	$\min(\{E\}) \hat{=} E$	where E is a single expression
*	SIMP_MAX_SING	$\max(\{E\}) \hat{=} E$	where E is a single expression
*	SIMP_MIN_NATURAL	$\min(\mathbb{N}) \hat{=} 0$	
*	SIMP_MIN_NATURAL1	$\min(\mathbb{N}_1) \hat{=} 1$	
*	SIMP_MIN_BUNION_SING	$\min(S \cup \dots \cup \{\min(T)\} \cup \dots \cup U) \hat{=} \min(S \cup \dots \cup T \cup \dots \cup U)$	
*	SIMP_MAX_BUNION_SING	$\max(S \cup \dots \cup \{\max(T)\} \cup \dots \cup U) \hat{=} \max(S \cup \dots \cup T \cup \dots \cup U)$	
*	SIMP_MIN_UPTO	$\min(E .. F) \hat{=} E$	
*	SIMP_MAX_UPTO	$\max(E .. F) \hat{=} F$	
*	SIMP_LIT_MIN	$\min(\{E, \dots, i, \dots, j, \dots, H\}) \hat{=} \min(\{E, \dots, i, \dots, H\})$	where i and j are literals and $i \leq j$
*	SIMP_LIT_MAX	$\max(\{E, \dots, i, \dots, j, \dots, H\}) \hat{=} \max(\{E, \dots, i, \dots, H\})$	where i and j are literals and $i \leq j$
*	SIMP_SPECIAL_CARD	$\text{card}(\emptyset) \hat{=} 0$	
*	SIMP_CARD_SING	$\text{card}(\{E\}) \hat{=} 1$	where E is a single expression
*	SIMP_SPECIAL_EQUAL_CARD	$\text{card}(S) = 0 \hat{=} S = \emptyset$	
*	SIMP_CARD_POW	$\text{card}(\mathbb{P}(S)) \hat{=} 2^{\text{card}(S)}$	
*	SIMP_CARD_BUNION	$\text{card}(S \cup T) \hat{=} \text{card}(S) + \text{card}(T) - \text{card}(S \cap T)$	
			with hypotheses

	SIMP_CARD_SETMINUS	$\text{card}(S \setminus T) \hat{=} \text{card}(S) - \text{card}(T)$	$T \subseteq S$ and either $\text{finite}(S)$ or $\text{finite}(T)$
	SIMP_CARD_SETMINUS_SETENUM	$\text{card}(S \setminus \{E_1, \dots, E_n\}) \hat{=} \text{card}(S) - \text{card}(\{E_1, \dots, E_n\})$	with hypotheses $E_i \in S$ for all $i \in 1 \dots n$
*	SIMP_CARD_CONVERSE	$\text{card}(r^{-1}) \hat{=} \text{card}(r)$	
*	SIMP_CARD_ID	$\text{card}(\text{id}) \hat{=} \text{card}(S)$	where id has type $\mathbb{P}(S \times S)$
*	SIMP_CARD_ID_DOMRES	$\text{card}(S \triangleleft \text{id}) \hat{=} \text{card}(S)$	
*	SIMP_CARD_PRJ1	$\text{card}(\text{prj}_1) \hat{=} \text{card}(S \times T)$	where prj_1 has type $\mathbb{P}(S \times T \times S)$
*	SIMP_CARD_PRJ2	$\text{card}(\text{prj}_2) \hat{=} \text{card}(S \times T)$	where prj_2 has type $\mathbb{P}(S \times T \times S)$
*	SIMP_CARD_PRJ1_DOMRES	$\text{card}(E \triangleleft \text{prj}_1) \hat{=} \text{card}(E)$	
*	SIMP_CARD_PRJ2_DOMRES	$\text{card}(E \triangleleft \text{prj}_2) \hat{=} \text{card}(E)$	
*	SIMP_CARD_LAMBDA	$\text{card}(\{x \cdot P \mid E \mapsto F\}) \hat{=} \text{card}(\{x \cdot P \mid E\})$	where E is a maplet combination of bound identifiers and expressions that are not bound by the comprehension (i.e., E is syntactically injective) and all identifiers bound in the comprehension set that occur in E also occur in F
*	SIMP_LIT_CARD_UPTO	$\text{card}(i .. j) \hat{=} j - i + 1$	where i and j are literals and $i \leq j$
	SIMP_TYPE_CARD	$\text{card}(T_{\text{enum}}) \hat{=} N$	where T_{enum} is a carrier set containing N elements
*	SIMP_LIT_GE_CARD_1	$\text{card}(S) \geq 1 \hat{=} \neg S = \emptyset$	
*	SIMP_LIT_LE_CARD_1	$1 \leq \text{card}(S) \hat{=} \neg S = \emptyset$	
*	SIMP_LIT_LE_CARD_0	$0 \leq \text{card}(S) \hat{=} \top$	
*	SIMP_LIT_GE_CARD_0	$\text{card}(S) \geq 0 \hat{=} \top$	
*	SIMP_LIT_GT_CARD_0	$\text{card}(S) > 0 \hat{=} \neg S = \emptyset$	
*	SIMP_LIT_LT_CARD_0	$0 < \text{card}(S) \hat{=} \neg S = \emptyset$	
*	SIMP_LIT_EQUAL_CARD_1	$\text{card}(S) = 1 \hat{=} \exists x \cdot S = \{x\}$	
*	SIMP_CARD_NATURAL	$\text{card}(S) \in \mathbb{N} \hat{=} \top$	
*	SIMP_CARD_NATURAL1	$\text{card}(S) \in \mathbb{N}_1 \hat{=} \neg S = \emptyset$	
*	SIMP_LIT_IN_NATURAL	$i \in \mathbb{N} \hat{=} \top$	where i is a non-negative literal
*	SIMP_SPECIAL_IN_NATURAL1	$0 \in \mathbb{N}_1 \hat{=} \perp$	
*	SIMP_LIT_IN_NATURAL1	$i \in \mathbb{N}_1 \hat{=} \top$	where i is a positive literal
*	SIMP_LIT_UPTO	$i .. j \hat{=} \emptyset$	where i and j are literals and $j < i$
*	SIMP_LIT_IN_MINUS_NATURAL	$-i \in \mathbb{N} \hat{=} \perp$	where i is a positive literal
*	SIMP_LIT_IN_MINUS_NATURAL1	$-i \in \mathbb{N}_1 \hat{=} \perp$	where i is a non-negative literal
*	DEF_IN_NATURAL	$x \in \mathbb{N} \hat{=} 0 \leq x$	

*	DEF_IN_NATURAL1	$x \in \mathbb{N}_1 \hat{=} 1 \leq x$	
*	SIMP_LIT_EQUAL_KBOOL_TRUE	$\text{bool}(P) = \text{TRUE} \hat{=} P$	
*	SIMP_LIT_EQUAL_KBOOL_FALSE	$\text{bool}(P) = \text{FALSE} \hat{=} \neg P$	
	DEF_EQUAL_MIN	$E = \min(S) \hat{=} E \in S \wedge (\forall x \cdot x \in S \Rightarrow E \leq x)$	where x non fre S, E
	DEF_EQUAL_MAX	$E = \max(S) \hat{=} E \in S \wedge (\forall x \cdot x \in S \Rightarrow E \geq x)$	where x non fre S, E
*	SIMP_SPECIAL_PLUS	$E + \dots + 0 + \dots + F \hat{=} E + \dots + F$	
*	SIMP_SPECIAL_MINUS_R	$E - 0 \hat{=} E$	
*	SIMP_SPECIAL_MINUS_L	$0 - E \hat{=} -E$	
*	SIMP_MINUS_MINUS	$-(-E) \hat{=} E$	
*	SIMP_MINUS_UNMINUS	$E - (-F) \hat{=} E + F$	where $(-F)$ is unary minus expression or a negative literal
*	SIMP_MULTI_MINUS	$E - E \hat{=} 0$	
*	SIMP_MULTI_MINUS_PLUS_L	$(A + \dots + C + \dots + B) - C \hat{=} A + \dots + B$	
*	SIMP_MULTI_MINUS_PLUS_R	$C - (A + \dots + C + \dots + B) \hat{=} -(A + \dots + B)$	
*	SIMP_MULTI_MINUS_PLUS_PLUS	$(A + \dots + E + \dots + B) - (C + \dots + E + \dots + D) \hat{=} (A + \dots + B) - (C + \dots + D)$	
*	SIMP_MULTI_PLUS_MINUS	$(A + \dots + D + \dots + (C - D) + \dots + B) \hat{=} A + \dots + C + \dots + B$	
*	SIMP_MULTI_ARITHREL_PLUS_PLUS	$A + \dots + E + \dots + B < C + \dots + E + \dots + D \hat{=} A + \dots + B < C + \dots + D$	where the root relation ($<$ here one of $\{=, <, \leq, >\}$,
*	SIMP_MULTI_ARITHREL_PLUS_R	$C < A + \dots + C + \dots + B \hat{=} 0 < A + \dots + B$	where the root relation ($<$ here one of $\{=, <, \leq, >\}$,
*	SIMP_MULTI_ARITHREL_PLUS_L	$A + \dots + C + \dots + B < C \hat{=} A + \dots + B < 0$	where the root relation ($<$ here one of $\{=, <, \leq, >\}$,
*	SIMP_MULTI_ARITHREL_MINUS_MINUS_R	$A - C < B - C \hat{=} A < B$	where the root relation ($<$ here one of $\{=, <, \leq, >\}$,
*	SIMP_MULTI_ARITHREL_MINUS_MINUS_L	$C - A < C - B \hat{=} B < A$	where the root relation ($<$ here one of $\{=, <, \leq, >\}$,
*	SIMP_SPECIAL_PROD_0	$E * \dots * 0 * \dots * F \hat{=} 0$	
*	SIMP_SPECIAL_PROD_1	$E * \dots * 1 * \dots * F \hat{=} E * \dots * F$	
*	SIMP_SPECIAL_PROD_MINUS_EVEN	$(-E) * \dots * (-F) \hat{=} E * \dots * F$	if an even numb of $-$
*	SIMP_SPECIAL_PROD_MINUS_ODD	$(-E) * \dots * (-F) \hat{=} -(E * \dots * F)$	if an odd numbe $-$
*	SIMP_LIT_MINUS	$-(i) \hat{=} (-i)$	where i is a liter
*	SIMP_LIT_EQUAL	$i = j \hat{=} \top \text{ or } \perp \text{ (computation)}$	where i and j ai literals
*	SIMP_LIT_LE	$i \leq j \hat{=} \top \text{ or } \perp \text{ (computation)}$	where i and j ai literals
*	SIMP_LIT_LT	$i < j \hat{=} \top \text{ or } \perp \text{ (computation)}$	where i and j ai literals
*	SIMP_LIT_GE	$i \geq j \hat{=} \top \text{ or } \perp \text{ (computation)}$	where i and j ai literals
*	SIMP_LIT_GT	$i > j \hat{=} \top \text{ or } \perp \text{ (computation)}$	where i and j ai literals

*	SIMP_DIV_MINUS	$(-E) \div (-F) \hat{=} E \div F$	
*	SIMP_SPECIAL_DIV_1	$E \div 1 \hat{=} E$	
*	SIMP_SPECIAL_DIV_0	$0 \div E \hat{=} 0$	
*	SIMP_SPECIAL_EXP_N_1_R	$E^1 \hat{=} E$	
*	SIMP_SPECIAL_EXP_N_1_L	$1^E \hat{=} 1$	
*	SIMP_SPECIAL_EXP_N_0	$E^0 \hat{=} 1$	
*	SIMP_MULTI_LE	$E \leq E \hat{=} \top$	
*	SIMP_MULTI_LT	$E < E \hat{=} \perp$	
*	SIMP_MULTI_GE	$E \geq E \hat{=} \top$	
*	SIMP_MULTI_GT	$E > E \hat{=} \perp$	
*	SIMP_MULTI_DIV	$E \div E \hat{=} 1$	
*	SIMP_MULTI_DIV_PROD	$(X * \dots * E * \dots * Y) \div E \hat{=} X * \dots * Y$	
*	SIMP_MULTI_MOD	$E \text{ mod } E \hat{=} 0$	
	DISTRIB_PROD_PLUS	$a * (b + c) \hat{=} (a * b) + (a * c)$	
	DISTRIB_PROD_MINUS	$a * (b - c) \hat{=} (a * b) - (a * c)$	
	DERIV_NOT_EQUAL	$\neg E = F \hat{=} E < F \vee E > F$	E and F must of Integer type

Extension Proof Rules

Rules that are marked with a * in the first column are implemented in the latest version of Rodin. Rules without a * are planned to be implemented in future versions. Other conventions used in these tables are described in The_Proving_Perspective_(Rodin_User_Manual)#Rewrite_Rules.

Rewrite Rules

	Name	Rule	Side Condition	A/M
*	SIMP_SPECIAL_COND_BTRUE	$\text{COND}(\top, E_1, E_2) \hat{=} E_1$		A
*	SIMP_SPECIAL_COND_BFALSE	$\text{COND}(\perp, E_1, E_2) \hat{=} E_2$		A
*	SIMP_MULTI_COND	$\text{COND}(C, E, E) \hat{=} E$		A

Inference Rules

	Name	Rule	Side Condition
*	DATATYPE_DISTINCT_CASE	$\frac{\mathbf{H}, x = c_1(p_{11}, \dots, p_{1k}) \vdash \mathbf{G} \quad \dots \quad \mathbf{H}, x = c_n(p_{n1}, \dots, p_{nl}) \vdash \mathbf{G}}{\mathbf{H} \vdash \mathbf{G}}$	where x has a datatype DT as type and appears free in \mathbf{G} ; DT has constructors c_1, \dots, c_n ; parameters p_{ij} are introduced as fresh identifiers
*	DATATYPE_INDUCTION	$\frac{\mathbf{H}, x = c_1(p_1, \dots, p_k), \mathbf{P}(p_{I_1}), \dots, \mathbf{P}(p_{I_l}) \vdash \mathbf{P}(x) \quad \dots}{\mathbf{H} \vdash \mathbf{P}(x)}$	where x has inductive datatype DT as type and appears free in \mathbf{P} ; $\{p_{I_1}, \dots, p_{I_l}\} \subseteq \{p_1, \dots, p_k\}$ are the inductive parameters (if any); an antecedent is created for every constructor c_i of DT ; all parameters are introduced as fresh identifiers; examples here

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