

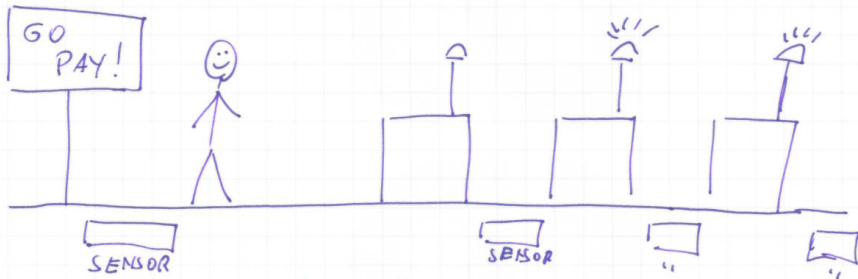
# A Checkout Counter Controller

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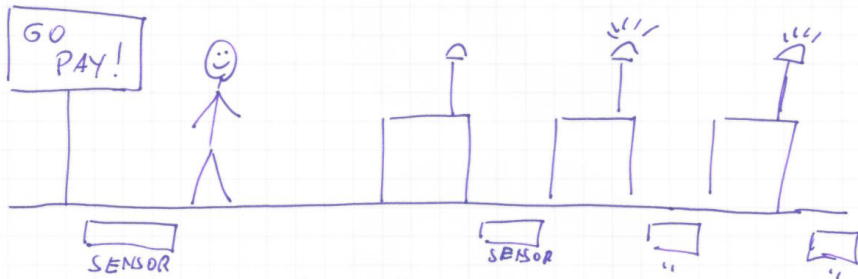
### Types of formulas

	Easy (few or no quantifiers)	More involved (quantifiers)
Sequential	Integer division, $n^2 = 1 + 3 + \dots$	Binary search
Concurrent, environment	Coffee Club	Checkout Counters



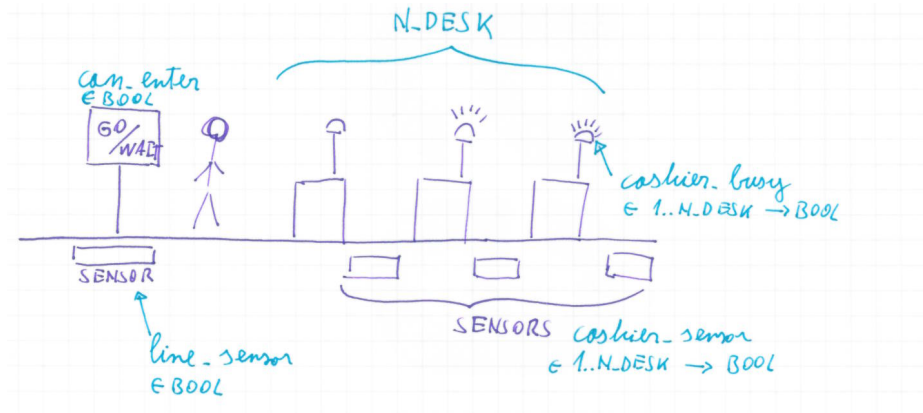
People **follow the rules**  
(otherwise automating is pointless or very difficult)

- Clients wait: entrance screen gives permission to go to counter.
- **Additionally:** clients wait for space between screen and counter to be empty.
- Clients do not leave counter before it has noticed client.



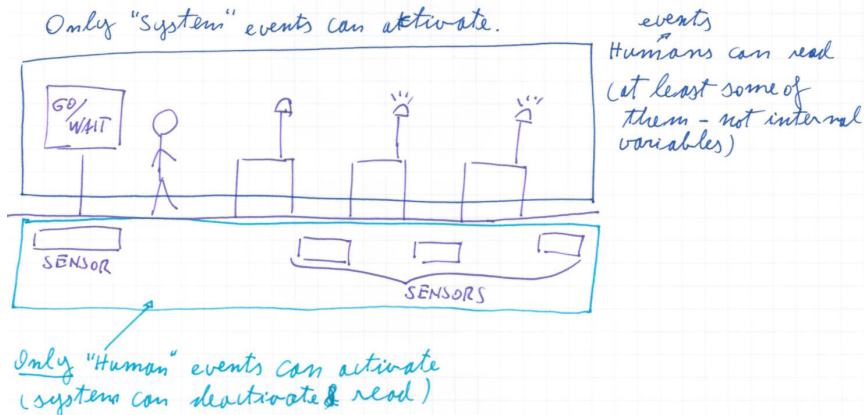
People **follow the rules**  
(otherwise automating is pointless or very difficult)

- Sensor after screen and in every counter: detect people.
- "Busy" light in every counter.
- People can go to any free counter.

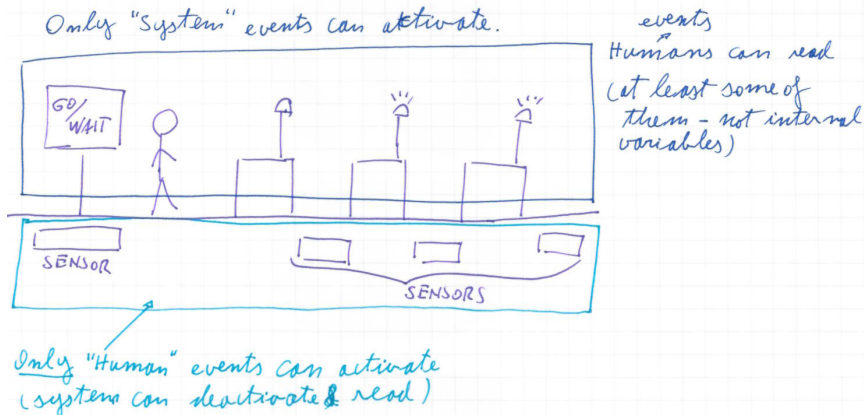


- “Signals” (screen, busy light, sensors, ...) modeled with variables.
- Plus internal state.
- Also system characteristic (e.g., # of counters).

# Types of Events



- Reactive system: neverending interaction loop (with humans).
- Correctness: we need to *simulate* (acceptable) human behavior.
- In general: model environment to ensure system works correctly.

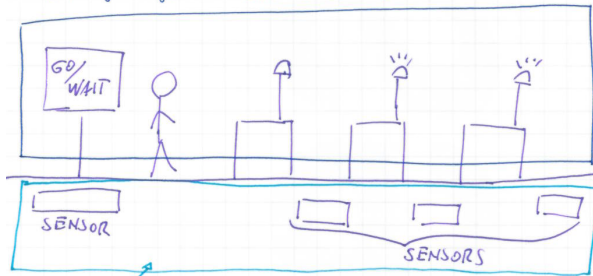


## Human Environment events

- Environment events read from some *system variables*.
- Give input to system by writing to variables modeling sensors.
- **Cannot write** on system variables / read from internal system variables.

# Types of Events

Only "System" events can activate.



events  
Humans can read  
(at least some of  
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Only "Human" events can activate  
(system can deactivate & read)

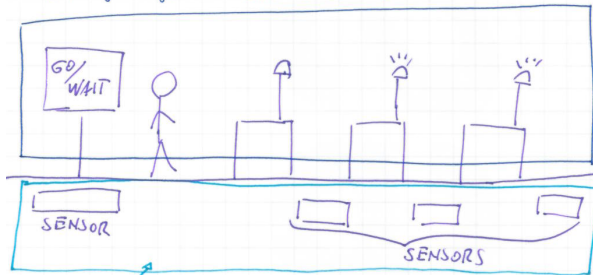
## System events

- Read and write system variables.
- Read environment variables.
- Cannot write onto environment variables.



# Types of Events

Only "System" events can activate.



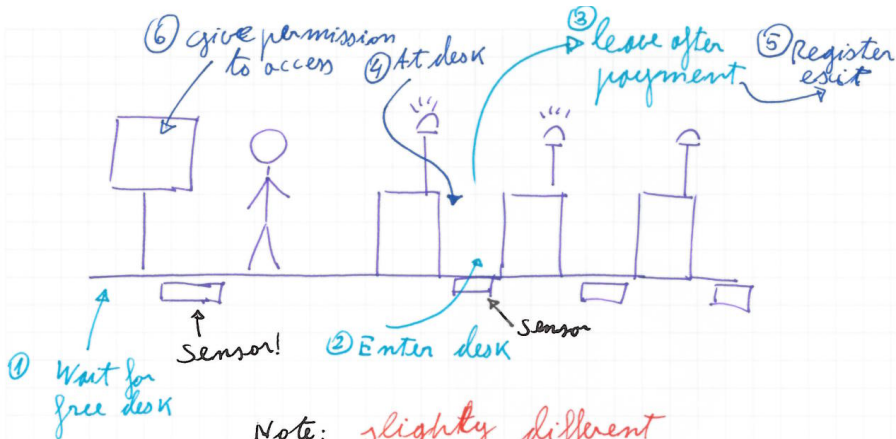
events  
Humans can read  
(at least some of  
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Only "Human" events can activate  
(system can deactivate & read)

- Environment events reading from internal system variables.
- Environment events writing to system variables.
- System events writing to environment variables altering environment behavior.

is cheating.

# Events in the Model



① Wait for free desk

3 Human

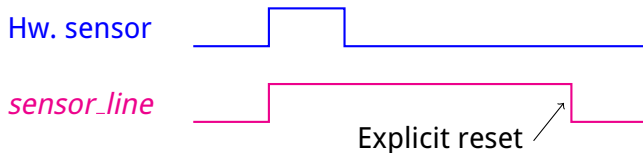
3 System

Note: slightly different behavior for sensors.

Can be simulated via software.

## Understanding Sensors

- Hardware usually on / off when detecting / not detecting.
- Transient behavior: real-time signals may be missed.  
(Unlikely for our case)
- We need to detect clients walking to counter (*sensor\_line*).
  - Sensor is *on* as long as clients walking to counter, or
  - Sensor *off* after clients pass, *sensor\_line* stays *on* until turned *off*.
    - The latter can be simulated with software.
    - Assume that behavior, do not show code.



- *cashier\_sensor(k)* is *on* only when client at counter.

## Making Sense of Sensor and Screen

*line\_sensor*

		False	True
<i>can_enter</i>	False	No free desk Corridor empty	No free desk Corridor not empty
	True	Free desk Corridor empty	Just entered

- All combinations possible.
- They summarize the immediate previous situation.
- Again, we assume **clients wait for corridor to be empty**.
- We could force it with a barrier:

$$\text{barrier} = \text{OPEN} \Leftrightarrow \text{can\_enter} = \text{TRUE} \wedge \text{line\_sensor} = \text{FALSE}$$

## Busy Lights and Counter Sensors

- Not in sync.
- That is how it physically happens.
- Every combination has a meaning: captures change over time.

		<i>busy(k)</i>	
		False	True
<i>sensor(k)</i>	False	Free	Just left
	True	Just arrived	Being served

## Some Invariants

If we can enter, there is a free desk

$$can\_enter = TRUE \Rightarrow \exists k \cdot (k \in 1..N\_DESK \wedge cashier\_busy(k) = FALSE)$$

Note: implication is not causality.

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If we are in the corridor and not yet served, there is a free desk

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Note: the two above need additional auxiliary invariants (see model).

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What about “entering should be disallowed if there is someone in the corridor”?

$$\neg (can\_enter = TRUE \wedge line\_sensor = TRUE)$$

Remember a previous slide: all combinations *can\_enter* and *line\_sensor* are legal and have a different meaning.



- Absolute time meaningless: cannot rely on relative speed.
- Safety, liveness.
- Auxiliary invariants (remember  $n^2 = 1 + 3 + \dots + (2n + 1)$ ).
- Function initialization.
- Deciding events: observables which change state. Be frugal!
- Modeling physical environment: very often a safe overapproximation.
- Add desk number in screen:
  - Type invariant:  $cashier\_number \in 0..N\_DESK$ .
  - Gluing invariant:  $cashier\_number = 0 \Leftrightarrow can\_enter = FALSE$ .
  - $cashier\_number = 0$  means “cannot enter”.
  - $cashier\_number : \in \{k | k \in 1..N\_DESK \wedge cashier\_busy(k) = FALSE \wedge cashier\_sensor(k) = FALSE\}$
  - Simple, straightforward refinement.